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Birational Rowmotion on a Rectangle over a
Noncommutative Ring

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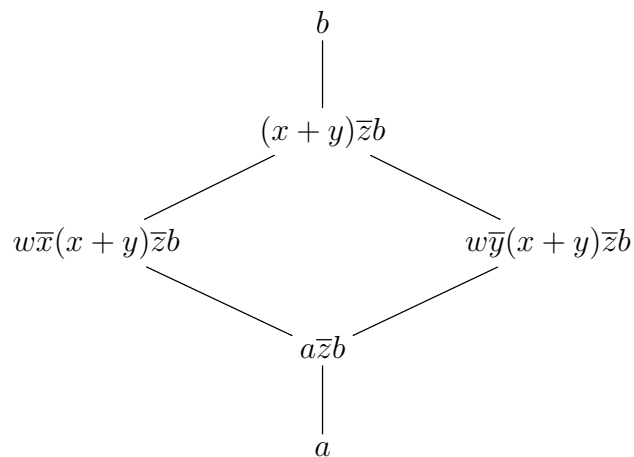
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Birational rowmotion on a rectangle over a noncommutative ring.

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Brief abstract. We extend the periodicity of birational rowmotion for rectangular posets to the case when the base field is replaced by a noncommutative ring (under appropriate conditions). This resolves a conjecture from 2014. The proof uses a novel approach and is fully self-contained.

Keywords: rowmotion; posets; noncommutative rings; semirings; Zamolodchikov periodicity; root systems; promotion; trees; graded posets; Grassmannian; tropicalization.

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Birational rowmotion on a rectangle over a noncommutative ring

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Abstract. We extend the periodicity of birational rowmotion for rectangular posets to the case when the base field is replaced by a noncommutative ring (under appropriate conditions). This resolves a conjecture from 2014. The proof uses a novel approach and is fully self-contained.

Consider labellings of a finite poset P by $|P| + 2$ elements of a ring \mathbb{K} : one label associated with each poset element and two constant labels for the added top and bottom elements in \hat{P} . *Birational rowmotion* is a partial map on such labellings. It was originally defined by Einstein and Propp [EinPro13] for $\mathbb{K} = \mathbb{R}$ as a lifting (via detropicalization) of *piecewise-linear rowmotion*, a map on the order polytope $\mathcal{O}(P) := \{\text{order-preserving } f : P \rightarrow [0, 1]\}$. The latter, in turn, extends the well-studied rowmotion map on the set of order ideals (or more properly, the set of order filters) of P , which correspond to the vertices of $\mathcal{O}(P)$. Dynamical properties of these combinatorial maps sometimes (but not always) extend to the birational level, while results proven at the birational level always imply their combinatorial counterparts. Allowing \mathbb{K} to be noncommutative, we generalize the birational level even further, and some properties are in fact lost at this step.

In 2014, the authors gave the first proof of periodicity for birational rowmotion on rectangular posets (when P is a product of two chains) for \mathbb{K} a field, and conjectured that it survives (in an appropriately twisted form) in

the noncommutative case. In this paper, we prove this noncommutative periodicity and a concomitant antipodal reciprocity formula. We end with some conjectures about periodicity for other posets, and the question of whether our results can be extended to (noncommutative) semirings.

It has been observed by Glick and Grinberg that, in the commutative case, periodicity of birational rowmotion can be used to derive Zamolodchikov periodicity in the type AA case, and vice-versa. However, for noncommutative \mathbb{K} , Zamolodchikov periodicity fails even in small examples (no matter what order the factors are multiplied), while noncommutative birational rowmotion continues to exhibit periodicity. Thus, our result can be viewed as a lateral generalization of Zamolodchikov periodicity to the noncommutative setting.

Keywords: rowmotion; posets; noncommutative rings; semirings; Zamolodchikov periodicity; root systems; promotion; trees; graded posets; Grassmannian; tropicalization.

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Introduction

The goal of this paper is to extend the periodicity of birational rowmotion for rectangular posets to the case when the base field is replaced by a noncommutative ring (under appropriate conditions). This resolves a conjecture from 2014. The proof uses a novel approach (even in the commutative case) and is fully self-contained.

Let P be a finite poset, and let \widehat{P} be the same poset with two extra elements added: one global minimum and one global maximum. For the time being, let \mathbb{K} be a field. A \mathbb{K} -labelling of P means a map from \widehat{P} to \mathbb{K} ; we view it as a way of labelling each element of \widehat{P} by an element of \mathbb{K} . Birational rowmotion, as studied conventionally, is a rational map R on such labellings (i.e., a rational map $R : \mathbb{K}^{\widehat{P}} \dashrightarrow \mathbb{K}^{\widehat{P}}$). It was introduced by Einstein and Propp [EinPro13] for $\mathbb{K} = \mathbb{R}$, generalizing (via the tropical limit) the well-studied *combinatorial rowmotion* map on order ideals of P [BrSchr74, StWi11, ProRob13, ThoWil19].

Birational rowmotion can be defined as a composition of “toggles”: For each $v \in P$, we define the v -toggle as the rational map $T_v : \mathbb{K}^{\widehat{P}} \dashrightarrow \mathbb{K}^{\widehat{P}}$ that modifies a \mathbb{K} -labelling f by changing the label $f(v)$ to¹

$$\left(\sum_{\substack{u \in \widehat{P}; \\ u \triangleleft v}} f(u) \right) \cdot (f(v))^{-1} \cdot \left(\sum_{\substack{u \in \widehat{P}; \\ u \triangleright v}} (f(u))^{-1} \right)^{-1},$$

while leaving all the other labels of f unchanged. Now, birational rowmotion R is the composition of all the v -toggles, where v runs over the poset P from top to bottom. (That is, we pick a linear extension (v_1, v_2, \dots, v_n) of P , and set $R = T_{v_1} \circ T_{v_2} \circ \dots \circ T_{v_n}$.)

¹The notations \triangleleft and \triangleright mean “covered by” and “covers”, respectively (see Sections 1 and 3 for details).

Dynamical properties at the combinatorial level sometimes extend to higher levels, while results proven at the birational level always imply their combinatorial counterparts. In particular, while combinatorial rowmotion always has finite order (since it is an invertible map on a finite set), there is no reason to expect periodicity at all at the higher levels. Indeed, for many nice posets, birational rowmotion has infinite order, including for the Boolean algebra of order 3 (or those in [Roby15, Fig. 6]), and there are only a few infinite classes where it appears to have finite order (mostly posets associated with representation theory, e.g., root or minuscule posets). In these cases the order of birational rowmotion is generally the *same* as for combinatorial rowmotion, e.g., $p + q$ for $P = [p] \times [q]$.

In 2014, the authors gave the first proof of periodicity of birational rowmotion for rectangular posets (i.e., when P is a product of two chains) and \mathbb{K} a field [GriRob14]. The main idea of this proof was to embed the space of labellings into an appropriate Grassmannian (where in each “sufficiently generic” \mathbb{K} -labelling, the labels can be expressed as ratios of certain minors of a matrix) and use particular Plücker relations to derive the result. There were several serious technical hurdles to overcome.

The definition of birational rowmotion relies entirely on addition, multiplication and inverses in \mathbb{K} . Thus, it is natural to extend it to the case when \mathbb{K} is a ring (not necessarily commutative), or even just a semiring. (At this level, birational rowmotion is no longer a rational map, just a partial map.) However, there is no guarantee that the properties of birational rowmotion survive at this level for every poset; and indeed, sometimes they do not (see, e.g., Example 13.9). However, in 2014, the authors experimentally observed that the periodicity for rectangular posets appears to hold even in this noncommutative setting, as long as it is appropriately modified: After $p + q$ iterations of birational rowmotion, the labels are not returned to their original states, but rather to certain “twisted variants” thereof (resembling, but not the same as, conjugates). See Example 3.17 to get the sense of this.

Strikingly, this noncommutative generalization has resisted all approaches that have previously succeeded in the commutative case. The determinantal computations involved in the proof in [GriRob14] can be extended to the noncommutative setting using the *quasideterminants* of Gelfand and Retakh, but it seems impossible to make a rigorous proof out of it (lacking, e.g., any useful notation of Zariski topology in this setting, it is not clear what it means for a \mathbb{K} -labelling to be “generic”). The alternative proof of commutative periodicity found by Musiker and Roby [MusRob17] (via a lattice-path formula for iterates of birational rowmotion) could not be generalized as well. Thus the noncommutative case remained an open problem.²

At some point, Glick and Grinberg noticed that the Y -variables in the type- AA Zamolod-

²This is not the first time that rational maps in algebraic combinatorics have been generalized to the noncommutative case; some other instances are [IyuShk14, BerRet15, Rupel17, GonKon21]. Each time, the generalizations have been much harder to prove, not least because very little of the commutative groundwork is (currently?) available at the noncommutative level. For instance, it is insufficient to work over the “free skew fields”, since an identity between rational expressions can be true in all skew fields yet fail in some noncommutative rings (such as the identity $x(yx)^{-1}y = 1$). For this reason, while natural from an algebraic point of view, the noncommutative setting is only recently and slowly getting explored.

chikov periodicity theorem of Volkov [Volk06] could be written as ratios of labels under iterated birational rowmotion [Roby15, § 4.4]; this allows the periodicity in one setting to be derived from that in the other (with some work). However, for noncommutative \mathbb{K} , Zamolodchikov periodicity fails even in small examples such as $r = r' = 2$ (no matter what order we multiply the factors), while noncommutative birational rowmotion continues to exhibit periodicity. This approach is therefore unavailable in the noncommutative case as well.

In this paper, we prove the periodicity of birational rowmotion and a concomitant antipodal reciprocity formula over an arbitrary noncommutative ring. The proof proceeds from first principles, by studying certain values V_ℓ^v and A_ℓ^v and their products along paths in the rectangle. An important step is the reduction of the claims to the labels on the “lower boundary” of the rectangle (i.e., to the labels at the elements of the form $(i, 1)$ and $(1, j)$). This reduction requires subtraction, which is why we are only addressing the case of a ring, not of a semiring; the latter remains open.

A few words are in order about the relation between our birational rowmotion and a parallel construction. Combinatorial rowmotion seems first to have been defined not on the set $J(P)$ of order ideals of P , but rather on the set $\mathcal{A}(P)$ of *antichains* of P [BrSchr74]. The standard bijection between $J(P)$ and $\mathcal{A}(P)$ (by taking maximal elements of $I \in J(P)$ or saturating down from an antichain) makes it easy to go between the two maps and to see that they have the same periodicity. However, some dynamic properties (e.g., homomesy) that depend on the sets themselves are not so easily translated. Just as Einstein and Propp lifted combinatorial rowmotion on $J(P)$ to a birational map and we continued to the noncommutative context, Joseph and Roby did a parallel lifting on the antichain side: from antichain rowmotion to piecewise-linear rowmotion on the *chain polytope*, $\mathcal{C}(P)$, to birational antichain rowmotion, and finally to noncommutative antichain rowmotion [JosRob20, JosRob21]. In particular they lifted “transfer maps” (originally defined by Stanley to go between $\mathcal{O}(P)$ and $\mathcal{C}(P)$ [Stan86]) from the piecewise-linear to the birational and noncommutative realms. These serve as equivariant bijections at each level, thus showing that periodicity at each level is equivalent for the order-ideal and antichain liftings. But they were unable to find a new proof of periodicity for the piecewise-linear and higher levels, relying instead on the periodicity results for birational order-ideal rowmotion to deduce it for birational antichain rowmotion. They also lifted a useful invariant, the *Stanley–Thomas word*, which cyclically rotates with antichain rowmotion at each level. At the combinatorial level, this gives an equivariant bijection that proves periodicity [ProRob13, § 3.3.2]; however, it is no longer a bijection at higher levels. Our paper completes the story in the case of a ring: Via the transfer maps mentioned above, the periodicity of noncommutative birational order-ideal rowmotion entails the periodicity of noncommutative birational antichain rowmotion.

The paper is structured in a fairly straightforward way: In the first sections (Sections 1 to 3), we introduce our noncommutative setup and define birational rowmotion in it. These include technicalities about partial maps and the definition of noncommutative toggles. In Section 4, we state our main results. In the sections that follow, we build an arsenal of lemmas to prove these results; the proofs are completed in Section 11. In the

final Sections 12 and 13, we discuss avenues for further work: a possible generalization to semirings and conjectured periodicity claims for other posets.

0.1. Remark on the level of detail

This paper comes in two versions: a regular one and a more detailed one. The regular version is optimized for readability, leaving out the more straightforward parts and technical arguments. The more detailed version has many of them expanded.

This is the regular version of the paper. The more detailed one can be obtained by replacing

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\excludecomment{verlong}
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\includecomment{vershort}
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by

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\includecomment{verlong}
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\excludecomment{vershort}
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in the preamble of the LaTeX sourcecode and then compiling to PDF. It is also available as an ancillary file on the arXiv page of this paper.

0.2. Acknowledgments

We are greatly indebted to the Mathematisches Forschungsinstitut Oberwolfach, which hosted us for three weeks during Summer 2021. Much of this paper was conceived during that stay. We thank Gerhard Huisken and Andrea Schillinger in particular for their flexibility in the scheduling of the visit.

We are also grateful to Banff International Research Station for hosting a hybrid workshop on dynamical algebraic combinatorics in November 2021 where these results were first presented.

We further acknowledge our appreciation of Michael Joseph, Tim Campion, Max Glick, Maxim Kontsevich, Gregg Musiker, Pace Nielsen, James Propp, Pasha Pylyavskyy, Bruce Sagan, Roland Speicher, David Speyer, Hugh Thomas, and Jurij Volcic, for useful advice and conversations.

Computations using the SageMath computer algebra system [S⁺09] provided essential data for us to conjecture some of the results.

1. Linear extensions of posets

This section collects a few standard notions concerning posets and their linear extensions, needed to define the main characters of our paper. Readers familiar with the subject may wish to skip forward to Section 2 or Section 3. We start by defining general notations identical with those in [GriRob14], to which we refer the reader for commentary and comparison to other references.

Convention 1.1. We let \mathbb{N} denote the set $\{0, 1, 2, \dots\}$.

Definition 1.2. Let P be a poset, and $u, v \in P$.

- (a) We will use the symbols \leq , $<$, \geq and $>$ to denote the lesser-or-equal relation, the lesser relation, the greater-or-equal relation and the greater relation, respectively, of the poset P . (Thus, for example, “ $u < v$ ” means “ u is smaller than v with respect to the partial order on P ”.)
- (b) The elements u and v of P are said to be *incomparable* if we have neither $u \leq v$ nor $u \geq v$.
- (c) We write $u \triangleleft v$ if we have $u < v$ and there is no $w \in P$ such that $u < w < v$. One often says that “ u is covered by v ” to signify that $u \triangleleft v$.
- (d) We write $u \triangleright v$ if we have $u > v$ and there is no $w \in P$ such that $u > w > v$. (Thus, $u \triangleright v$ holds if and only if $v \triangleleft u$.) One often says that “ u covers v ” to signify that $u \triangleright v$.
- (e) An element u of P is called *maximal* if every $w \in P$ satisfying $w \geq u$ satisfies $w = u$. In other words, an element u of P is called *maximal* if there is no $w \in P$ such that $w > u$.
- (f) An element u of P is called *minimal* if every $w \in P$ satisfying $w \leq u$ satisfies $w = u$. In other words, an element u of P is called *minimal* if there is no $w \in P$ such that $w < u$.

These notations may become ambiguous when an element belongs to several different posets simultaneously. In such cases, we will disambiguate them by adding the words “in P ” (where P is the poset which we want to use).³

Convention 1.3. From now on, for the rest of the paper, **we fix a finite poset P** . Most of our results will concern the case when P has a rather specific form (viz., a rectangular poset, i.e., a Cartesian product of two finite chains), but we do not assume this straightaway.

Definition 1.4. A *linear extension* of P will mean a list (v_1, v_2, \dots, v_m) of the elements of P such that

- each element of P occurs exactly once in this list, and
- any $i, j \in \{1, 2, \dots, m\}$ satisfying $v_i < v_j$ (in P) must satisfy $i < j$ (in \mathbb{Z}).

³For instance, if R denotes the poset \mathbb{Z} endowed with the **reverse** of its usual order, then we say (for instance) that “ $0 > 3$ in R ” rather than just “ $0 > 3$ ” (to avoid mistaking our statement for an absurd claim about the usual order on \mathbb{Z}).

A linear extension of P is also known as a *topological sorting* of P . We will use the following well-known fact:

■ **Theorem 1.5.** There exists a linear extension of P .

■ **Definition 1.6.** The set of all linear extensions of P will be called $\mathcal{L}(P)$. Thus, $\mathcal{L}(P) \neq \emptyset$ (by Theorem 1.5).

The reader can easily verify the following proposition:

■ **Proposition 1.7.** Let (v_1, v_2, \dots, v_m) be a linear extension of P . Let $i \in \{1, 2, \dots, m-1\}$ be such that the elements v_i and v_{i+1} of P are incomparable. Then $(v_1, v_2, \dots, v_{i-1}, v_{i+1}, v_i, v_{i+2}, v_{i+3}, \dots, v_m)$ (this is the tuple obtained from the tuple (v_1, v_2, \dots, v_m) by interchanging the adjacent entries v_i and v_{i+1}) is a linear extension of P as well.

We will also use the following folklore result:⁴

■ **Proposition 1.8.** Let \sim denote the equivalence relation on $\mathcal{L}(P)$ generated by the following requirement: For any linear extension (v_1, v_2, \dots, v_m) of P and any $i \in \{1, 2, \dots, m-1\}$ such that the elements v_i and v_{i+1} of P are incomparable, we set

$$(v_1, v_2, \dots, v_m) \sim (v_1, v_2, \dots, v_{i-1}, v_{i+1}, v_i, v_{i+2}, v_{i+3}, \dots, v_m).$$

Then any two elements of $\mathcal{L}(P)$ are equivalent under the relation \sim .

Proofs of Proposition 1.8 can be found in [GriRob14, Proposition 1.7], in [AyKlSc12, Proposition 4.1 (for the $\pi' = \pi\tau_j$ case)] and in [Etienn84, Lemma 1]. See also [Naatz00, Proposition 2.2] for a stronger claim (describing a shortest way to transform a given linear extension into another by successively swapping adjacent incomparable entries).

Another well-known fact says that any nonempty finite poset has a minimal element and a maximal element. In other words:

■ **Proposition 1.9.** Assume that $P \neq \emptyset$. Then:

- (a) The poset P has a minimal element.
- (b) The poset P has a maximal element.

2. Inverses in rings

⁴Particular cases of Proposition 1.8 have a tendency to appear in various parts of combinatorics; see [DefKra21, Proposition 1.3] for a few such references.

Convention 2.1. From now on, for the rest of this paper, **we fix a ring** \mathbb{K} . This ring is not required to be commutative, but must have a unity and be associative.

For example, \mathbb{K} can be \mathbb{Z} or \mathbb{Q} or \mathbb{C} or a polynomial ring or a matrix ring over any of these. In almost all previous work on birational rowmotion (with the exception of [JosRob20] and [JosRob21]), only commutative rings (and, occasionally, semirings) were considered; by removing the commutativity assumption, we are invalidating many of the methods used in prior research. We suspect that the level of generality can be increased even further, replacing our ring \mathbb{K} by a semiring (i.e., a “ring without subtraction”); however, this poses new difficulties which we will not surmount in the present work. (See Section 12 for more about this.)

Even as we do not assume our ring \mathbb{K} to be a division ring, we will nevertheless take multiplicative inverses of elements of \mathbb{K} on many occasions. These inverses do not always exist, but when they do exist, they are unique; thus, we introduce a notation for them:

Definition 2.2. Let a be an element of \mathbb{K} .

- (a) An *inverse* of a means an element $b \in \mathbb{K}$ such that $ab = ba = 1$. This inverse is unique when it exists, and will be denoted by \bar{a} . (A more standard notation for it is a^{-1} , but we prefer the notation \bar{a} since it helps keep our formulas short.)
- (b) We say that the element a of \mathbb{K} is *invertible* if it has an inverse.

The following well-known properties of inverses will often be used without mention:

Proposition 2.3.

- (a) If a is an invertible element of \mathbb{K} , then its inverse \bar{a} is invertible as well, and its inverse is $\overline{\bar{a}} = a$.
- (b) If a and b are two invertible elements of \mathbb{K} , then their product ab is invertible as well, and its inverse is $\overline{ab} = \bar{b} \cdot \bar{a}$.
- (c) If a_1, a_2, \dots, a_m are several invertible elements of \mathbb{K} , then their product $a_1 a_2 \cdots a_m$ is invertible as well, and its inverse is $\overline{a_1 a_2 \cdots a_m} = \bar{a}_m \cdot \bar{a}_{m-1} \cdots \bar{a}_1$.

The converse of Proposition 2.3 (b) does not necessarily hold: A product ab of two elements a and b of \mathbb{K} can be invertible even when neither a nor b is⁵.

The next property of inverses is less well-known:⁶

⁵See <https://math.stackexchange.com/questions/627562> for examples of such situations.

⁶Proposition 2.4 (a) will not be used in what follows, but its proof provides a good warm-up exercise in manipulating inverses in a noncommutative ring.

Proposition 2.4. Let a and b be two elements of \mathbb{K} such that $a+b$ is invertible. Then:

(a) We have $a \cdot \overline{a+b} \cdot b = b \cdot \overline{a+b} \cdot a$.

(b) If both a and b are invertible, then $\bar{a} + \bar{b}$ is invertible as well and its inverse is

$$\overline{\bar{a} + \bar{b}} = a \cdot \overline{a+b} \cdot b.$$

Proof. (a) Comparing

$$a \cdot \overline{a+b} \cdot a + a \cdot \overline{a+b} \cdot b = a \cdot \underbrace{\overline{a+b} \cdot (a+b)}_{=1} = a$$

with

$$a \cdot \overline{a+b} \cdot a + b \cdot \overline{a+b} \cdot a = \underbrace{(a+b) \cdot \overline{a+b}}_{=1} \cdot a = a,$$

we obtain

$$a \cdot \overline{a+b} \cdot a + a \cdot \overline{a+b} \cdot b = a \cdot \overline{a+b} \cdot a + b \cdot \overline{a+b} \cdot a.$$

Subtracting $a \cdot \overline{a+b} \cdot a$ from both sides of this equality, we obtain $a \cdot \overline{a+b} \cdot b = b \cdot \overline{a+b} \cdot a$. This proves Proposition 2.4 (a).

(b) Assume that both a and b are invertible. Set $x := \bar{a} + \bar{b}$ and $y := a \cdot \overline{a+b} \cdot b$.

From $x = \bar{a} + \bar{b}$, we obtain $x \cdot a = (\bar{a} + \bar{b}) \cdot a = \underbrace{\bar{a}a}_{=1} + \bar{b}a = 1 + \bar{b}a$. Comparing this with

$$\bar{b} \cdot (a+b) = \bar{b}a + \underbrace{\bar{b}b}_{=1} = \bar{b}a + 1 = 1 + \bar{b}a,$$

we obtain $x \cdot a = \bar{b} \cdot (a+b)$. Now, from $y = a \cdot \overline{a+b} \cdot b$, we obtain

$$x \cdot y = \underbrace{x \cdot a}_{=\bar{b} \cdot (a+b)} \cdot \overline{a+b} \cdot b = \bar{b} \cdot \underbrace{(a+b) \cdot \overline{a+b}}_{=1} \cdot b = \bar{b} \cdot b = 1.$$

A similar argument (starting with $b \cdot x = b\bar{a} + 1 = (a+b) \cdot \bar{a}$) shows that $y \cdot x = 1$, so that y is an inverse of x . Hence, x is invertible and its inverse is $\bar{x} = y$. This is precisely the claim of Proposition 2.4 (b). \square

3. Noncommutative birational rowmotion

In this section, we introduce the basic objects whose nature we will investigate: labellings of a finite poset P by elements of a ring, and a partial map between them called “birational rowmotion”. These labellings generalize the field-valued labellings studied in [GriRob14], which in turn generalize the piecewise-linear labellings of [EinPro13], which in turn generalize the order ideals of P . Many of the definitions that follow will imitate analogous definitions made (in somewhat lesser generality) in [GriRob14].

3.1. The extended poset \widehat{P}

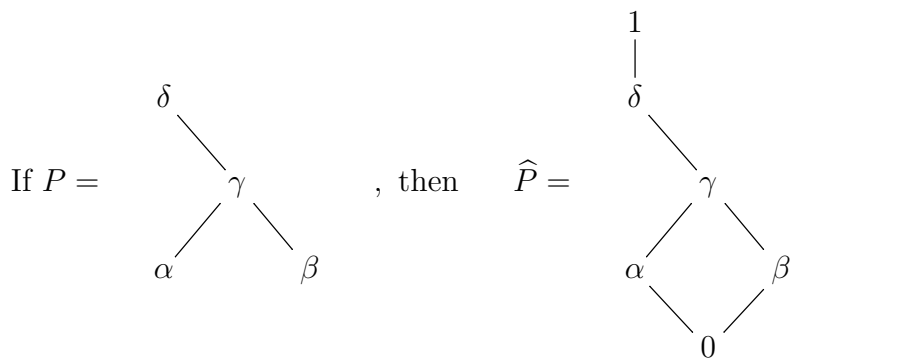
Definition 3.1. We define a poset \widehat{P} as follows: As a set, let \widehat{P} be the disjoint union of the set P with the two-element set $\{0, 1\}$. The smaller-or-equal relation \leq on \widehat{P} will be given by

$$(a \leq b) \iff ((a \in P \text{ and } b \in P \text{ and } a \leq b \text{ in } P) \text{ or } a = 0 \text{ or } b = 1).$$

Here and in the following, we regard the canonical injection of the set P into the disjoint union \widehat{P} as an inclusion; thus, P becomes a subset of \widehat{P} .

In the terminology of Stanley's [Stan11, Section 3.2], this poset \widehat{P} is the ordinal sum $\{0\} \oplus P \oplus \{1\}$.

Example 3.2. Let us represent posets by their Hasse diagrams. Then:



3.2. \mathbb{K} -labellings

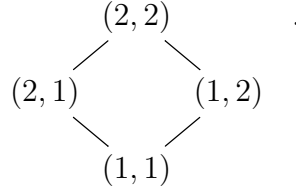
Let us now define the type of object on which our maps will act:

Definition 3.3. A \mathbb{K} -labelling of P will mean a map $f : \widehat{P} \rightarrow \mathbb{K}$. Thus, $\mathbb{K}^{\widehat{P}}$ is the set of all \mathbb{K} -labellings of P . If f is a \mathbb{K} -labelling of P and v is an element of \widehat{P} , then $f(v)$ will be called the *label of f at v* .

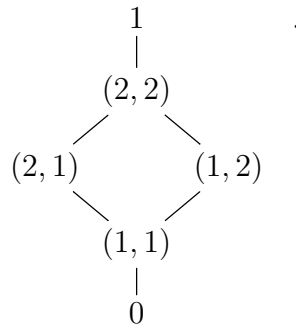
Example 3.4. Assume that P is the poset $\{1, 2\} \times \{1, 2\}$ with order relation defined by setting

$$(i, k) \leq (i', k') \quad \text{if and only if } (i \leq i' \text{ and } k \leq k').$$

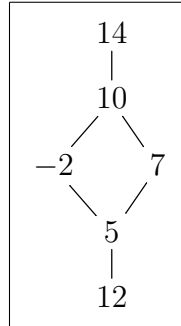
This poset will later be called the “ 2×2 -rectangle” in Definition 4.2. It has Hasse diagram



The extended poset \widehat{P} has Hasse diagram



We recall that a \mathbb{K} -labelling of P is a map $f : \widehat{P} \rightarrow \mathbb{K}$. We can visualize such a \mathbb{K} -labelling by replacing, in the Hasse diagram of \widehat{P} , each element $v \in \widehat{P}$ by the label $f(v)$. For example, the \mathbb{Z} -labelling of P that sends 0, $(1, 1)$, $(1, 2)$, $(2, 1)$, $(2, 2)$, and 1 to 12, 5, 7, -2 , 10, and 14, respectively can be visualized as follows:



For example, its label at $(1, 2)$ is 7.

3.3. Partial maps

We will next define the notion of a partial map, to formalize the idea of an operation whose result may be undefined, such as division on \mathbb{Q} (since division by zero is undefined). We

will use \perp as a symbol for such undefined values:

Convention 3.5. We fix an object called \perp . In the following, we tacitly assume that none of the sets we will consider contains this object \perp (unless otherwise specified).

The reader can think of \perp as a “division-by-zero error” (more precisely, a “division-by-a-non-invertible-element error”, since 0 is often not the only non-invertible element of \mathbb{K}).

Definition 3.6. Let X and Y be two sets. A *partial map* from X to Y means a map from X to $Y \sqcup \{\perp\}$.

If f is a partial map from X to Y , then f can be canonically extended to a map from $X \sqcup \{\perp\}$ to $Y \sqcup \{\perp\}$ by setting $f(\perp) := \perp$. We always consider f to be extended in this way.

If f is a partial map from X to Y , then the set $\{x \in X \mid f(x) \neq \perp\}$ will be called the *domain of definition* of f .

We view the element \perp as an “undefined output” – i.e., we think of a partial map f from X to Y as a “map” from X to Y that is defined only on some elements of X (namely, on those whose image under this map is not \perp). Thus, for example, in \mathbb{Q} , division is a partial map because division by 0 is undefined:

Example 3.7. The map

$$\begin{aligned} \mathbb{Q} &\rightarrow \mathbb{Q} \sqcup \{\perp\}, \\ x &\mapsto \begin{cases} 1/x, & \text{if } x \neq 0; \\ \perp, & \text{if } x = 0 \end{cases} \end{aligned}$$

is a partial map from \mathbb{Q} to \mathbb{Q} .

Partial maps can be composed much like usual maps:

Definition 3.8.

(a) Let X , Y and Z be three sets. Let f be a partial map from Y to Z . Let g be a partial map from X to Y .

Then $f \circ g$ denotes the partial map from X to Z that sends

$$\text{each } x \in X \text{ to } \begin{cases} f(g(x)), & \text{if } g(x) \neq \perp; \\ \perp, & \text{if } g(x) = \perp. \end{cases}$$

(Following our convention that $f(\perp)$ is understood to be \perp , we could simplify the right hand side to just $f(g(x))$, but we subdivided it into two cases just to stress the different “control flows”.)

This partial map $f \circ g$ is called the *composition* of f and g .

- (b) This notion of composition lets us define a category whose objects are sets and whose morphisms are partial maps. (The identity maps in this category are the obvious ones: i.e., the maps $\text{id} : X \rightarrow X \sqcup \{\perp\}$ that send each $x \in X$ to $x \in X \subseteq X \sqcup \{\perp\}$.)
- (c) Thus, if X is any set, and if f is any partial map from X to X , then we can define $f^k := \underbrace{f \circ f \circ \cdots \circ f}_{k \text{ times}}$ for any $k \in \mathbb{N}$.

Convention 3.9. Let X and Y be two sets. We will write “ $f : X \dashrightarrow Y$ ” for “ f is a partial map from X to Y ” (just as maps from X to Y are denoted “ $f : X \rightarrow Y$ ”).

A warning is in place: While we are using the symbol \dashrightarrow for partial maps here, the same symbol has been used for rational maps in [GriRob14]. The two uses are, in a sense, serving the same purpose (viz., describing “maps defined only on those inputs for which the relevant denominators are invertible”), but they have some technical differences. Rational maps are defined only when \mathbb{K} is an infinite field⁷, but are well-behaved in many ways that partial maps are not. (For example, a rational map is uniquely determined if its values on a Zariski-dense subset of its domain are known, but no such claims can be made for partial maps.) Thus, by working with partial maps instead of rational maps, we are freeing ourselves from technical assumptions on \mathbb{K} , but at the same time forcing ourselves to be explicit about the domains on which our partial maps are defined.

3.4. Toggles

Recall that $\mathbb{K}^{\hat{P}}$ denotes the set of \mathbb{K} -labellings of a poset P (that is, the set of all maps $\hat{P} \rightarrow \mathbb{K}$). Next, we define (*noncommutative*) *toggles*: a simple kind of partial maps on this set.

Definition 3.10. Let $v \in P$. We define a partial map $T_v : \mathbb{K}^{\hat{P}} \dashrightarrow \mathbb{K}^{\hat{P}}$ as follows: If $f \in \mathbb{K}^{\hat{P}}$ is any \mathbb{K} -labelling of P , then the \mathbb{K} -labelling $T_v f \in \mathbb{K}^{\hat{P}}$ is given by

$$(T_v f)(w) = \begin{cases} f(w), & \text{if } w \neq v; \\ \left(\sum_{\substack{u \in \hat{P}; \\ u < v}} f(u) \right) \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \hat{P}; \\ u > v}} f(u)}, & \text{if } w = v \end{cases} \quad (1)$$

for all $w \in \hat{P}$.

⁷It stands to reason that a notion of “rational map” should exist for a sufficiently wide class of infinite skew-fields as well, but we have not encountered a satisfactory theory of such maps in the literature. See <https://mathoverflow.net/questions/362724/> for a discussion of how this theory might start. It appears unlikely, however, that such “noncommutative rational maps” exist in the generality that we are working in (viz., arbitrary rings).

Here, we agree that if any part of the expression $\left(\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u) \right) \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > v}} f(u)}$ is not well-defined (i.e., if one of the values $f(u)$ and $f(v)$ appearing in it is undefined, or if $f(v)$ is not invertible, or if $f(u)$ is not invertible for some $u \in \widehat{P}$ satisfying $u > v$, or if the sum $\sum_{\substack{u \in \widehat{P}; \\ u > v}} \overline{f(u)}$ is not invertible), then $T_v f$ is understood to be \perp .

This partial map T_v is called the v -toggle or the toggle at v .

Thus, the partial map T_v is a “local” transformation: it only changes the label at the element v (unless its result is \perp).

Remark 3.11. You are reading Definition 3.10 right: We set $T_v f = \perp$ if any of $\overline{f(v)}$ and $\sum_{\substack{u \in \widehat{P}; \\ u > v}} \overline{f(u)}$ fails to be well-defined. Thus, in this case, none of the values $(T_v f)(w)$ exists. It may appear more natural to leave only the value $(T_v f)(v)$ undefined, while letting all other values $(T_v f)(w)$ equal the respective values $f(w)$. Our choice to “panic and crash”, however, will be more convenient for some of our proofs.

The v -toggle T_v is called a “noncommutative order toggle” in [JosRob20, Definition 5.6]. When the ring \mathbb{K} is commutative, this v -toggle T_v is an “involution” in the sense that each \mathbb{K} -labelling $f \in \mathbb{K}^{\widehat{P}}$ satisfying $T_v(T_v f) \neq \perp$ satisfies $T_v(T_v f) = f$. For noncommutative \mathbb{K} , this is usually not the case; an “inverse” partial map⁸ can be obtained by flipping the order of the factors on the right hand side of (1). (This “inverse” appears in [JosRob20] under the name “noncommutative order elgot”.)

The following proposition is trivially obtained by rewriting (1); we are merely stating it for easier reference in proofs:

Proposition 3.12. Let $v \in P$. For every $f \in \mathbb{K}^{\widehat{P}}$ satisfying $T_v f \neq \perp$, we have the following:

- (a) Every $w \in \widehat{P}$ such that $w \neq v$ satisfies $(T_v f)(w) = f(w)$.
- (b) We have

$$(T_v f)(v) = \left(\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u) \right) \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > v}} f(u)}.$$

⁸We are putting the word “inverse” in scare quotes since we are talking about partial maps, but the two maps are as close to being mutually inverse as partial maps can be.

Furthermore, the following “locality principle” (part of [JosRob20, Proposition 5.8]) is easy to check:⁹

Proposition 3.13. Let $v \in P$ and $w \in P$. Then $T_v \circ T_w = T_w \circ T_v$, unless we have either $v \triangleleft w$ or $w \triangleleft v$.

Proof of Proposition 3.13. In the case when \mathbb{K} is commutative, this is essentially [GriRob16, Proposition 14], except that we are now more careful about well-definedness (since only invertible elements have inverses). Yet, the proof given in [GriRob16] can easily be adapted to the general (noncommutative) case. The details can be found in the detailed version of this paper (but the reader should have an easy time reconstructing them). \square

As a particular case of Proposition 3.13, we have the following:

Corollary 3.14. Let v and w be two elements of P which are incomparable. Then $T_v \circ T_w = T_w \circ T_v$.

Corollary 3.15. Let (v_1, v_2, \dots, v_m) be a linear extension of P . Then the partial map $T_{v_1} \circ T_{v_2} \circ \dots \circ T_{v_m} : \mathbb{K}^{\hat{P}} \dashrightarrow \mathbb{K}^{\hat{P}}$ is independent of the choice of the linear extension (v_1, v_2, \dots, v_m) .

Proof. Combine Corollary 3.14 with Proposition 1.8. \square

3.5. Birational rowmotion

Recall that P is a finite poset. Corollary 3.15 lets us make the following definition.

Definition 3.16. *Birational rowmotion* (or, more precisely, the *birational rowmotion of P*) is defined as the partial map $T_{v_1} \circ T_{v_2} \circ \dots \circ T_{v_m} : \mathbb{K}^{\hat{P}} \dashrightarrow \mathbb{K}^{\hat{P}}$, where (v_1, v_2, \dots, v_m) is a linear extension of P . This partial map is well-defined, because

- Theorem 1.5 shows that a linear extension of P exists, and
- Corollary 3.15 shows that the partial map $T_{v_1} \circ T_{v_2} \circ \dots \circ T_{v_m}$ is independent of the choice of the linear extension (v_1, v_2, \dots, v_m) .

This partial map will be denoted by R .

Birational rowmotion is called “birational NOR-motion” (and denoted NOR) in [JosRob20, Definition 5.9]¹⁰. When \mathbb{K} is commutative, it agrees with the standard concept of birational rowmotion as studied in [EinPro13] and [GriRob14].

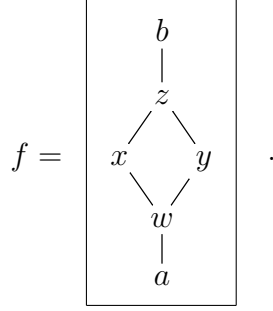
⁹In the following, equalities between partial maps are understood in the strongest possible sense: Two partial maps $F : X \dashrightarrow Y$ and $G : X \dashrightarrow Y$ satisfy $F = G$ if and only if each $x \in X$ satisfies $F(x) = G(x)$. This entails, in particular, that $F(x) = \perp$ holds if and only if $G(x) = \perp$. Thus, $F = G$ is a stronger requirement than merely saying that “ $F(x) = G(x)$ whenever neither $F(x)$ nor $G(x)$ is \perp ”.

¹⁰To be more precise, [JosRob20, Definition 5.9] works in a slightly less general context, requiring \mathbb{K} to be a skew field and that $f(0) = 1$ and $f(1) = C$ for some C in the center of \mathbb{K} .

Example 3.17. Let us demonstrate the effect of birational toggles and birational rowmotion. Namely, for this example, we let P be the poset $P = \{1, 2\} \times \{1, 2\}$ introduced in Example 3.4.

In order to disencumber our formulas, we agree to write $g(i, j)$ for $g((i, j))$ when g is a \mathbb{K} -labelling of P and (i, j) is an element of P .

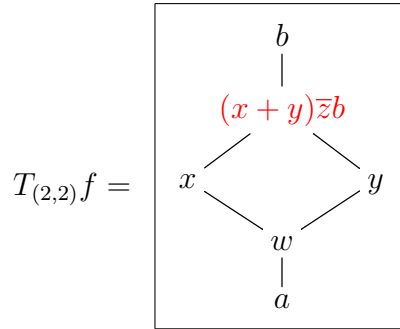
As in Example 3.4, we visualize a \mathbb{K} -labelling f of P by replacing, in the Hasse diagram of \widehat{P} , each element $v \in \widehat{P}$ by the label $f(v)$. Let f be a \mathbb{K} -labelling sending $0, (1, 1), (1, 2), (2, 1), (2, 2)$, and 1 to a, w, y, x, z , and b , respectively (for some elements a, b, x, y, z, w of \mathbb{K}); this f is then visualized as follows:



(As before, we draw $(2, 1)$ on the western corner and $(1, 2)$ on the eastern corner.)

Now, recall the definition of birational rowmotion R on our poset P . Since $((1, 1), (1, 2), (2, 1), (2, 2))$ is a linear extension of P , we have $R = T_{(1,1)} \circ T_{(1,2)} \circ T_{(2,1)} \circ T_{(2,2)}$. Let us track how this transforms our labelling f :

We first apply $T_{(2,2)}$, obtaining

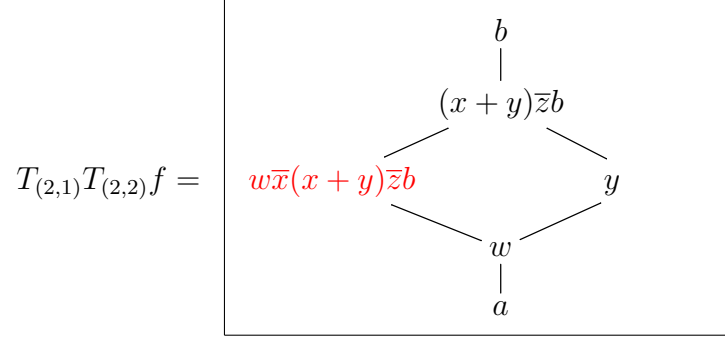


(where we colored the label at $(2, 2)$ red to signify that it is the label at the element which got toggled). Indeed, the only label that changes under $T_{(2,2)}$ is the one at $(2, 2)$, and this label becomes

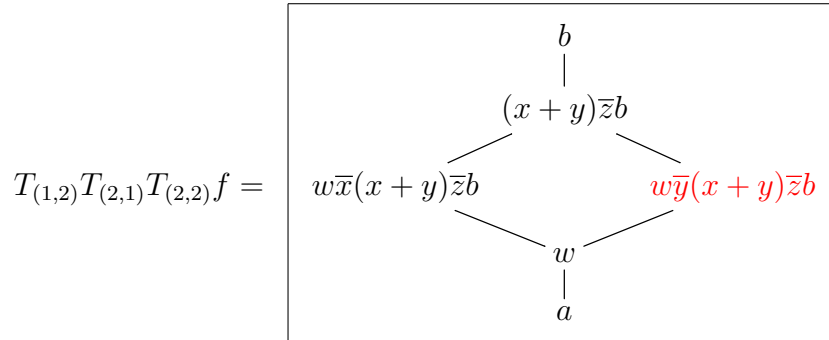
$$\begin{aligned} (T_{(2,2)}f)(2, 2) &= \left(\sum_{\substack{u \in \widehat{P}; \\ u < (2,2)}} f(u) \right) \cdot \overline{f(2, 2)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > (2,2)}} f(u)} \\ &= (f(1, 2) + f(2, 1)) \cdot \overline{f(2, 2)} \cdot \overline{f(1)} \\ &= (y + x) \cdot \bar{z} \cdot \bar{b} = (x + y) \cdot \bar{z} \cdot b. \end{aligned}$$

(We assume that z and b are indeed invertible; otherwise, $T_{(2,2)}f$ would be \perp and would remain \perp after any further toggles. Likewise, as we apply further toggles, we assume that everything else we need to invert is invertible.)

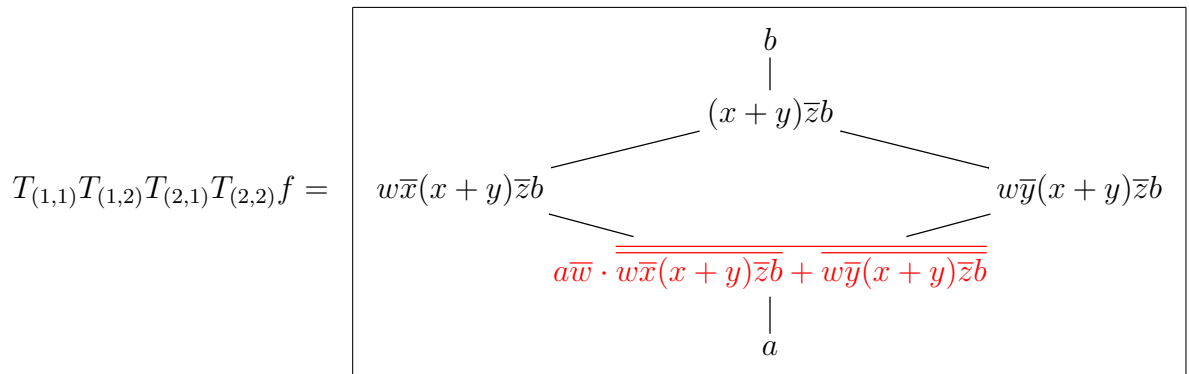
Having applied $T_{(2,2)}$, we next apply $T_{(2,1)}$, obtaining



Next, we apply $T_{(1,2)}$, obtaining



Finally, we apply $T_{(1,1)}$, resulting in



The unwieldy expression $\overline{\overline{a\bar{w} \cdot w\bar{x}(x+y)\bar{z}b + w\bar{y}(x+y)\bar{z}b}}$ in the label at (1,1) can be simplified to \bar{z} (using standard laws such as $\bar{p} \cdot \bar{q} = \overline{\overline{qp}}$ and distributivity), so this

rewrites as

$$T_{(1,1)}T_{(1,2)}T_{(2,1)}T_{(2,2)}f = \begin{array}{c} \boxed{\begin{array}{c} b \\ | \\ (x+y)\bar{z}b \\ / \quad \backslash \\ w\bar{x}(x+y)\bar{z}b \quad w\bar{y}(x+y)\bar{z}b \\ \backslash \quad / \\ a\bar{z}b \\ | \\ a \end{array}} \end{array} .$$

We thus have computed Rf (since $Rf = T_{(1,1)}T_{(1,2)}T_{(2,1)}T_{(2,2)}f$).

By repeating this procedure (or just substituting the labels of Rf obtained as variables), we can compute R^2f , R^3f etc., obtaining

$$Rf = \begin{array}{c} \boxed{\begin{array}{c} b \\ | \\ (x+y)\bar{z}b \\ / \quad \backslash \\ w\bar{x}(x+y)\bar{z}b \quad w\bar{y}(x+y)\bar{z}b \\ \backslash \quad / \\ a\bar{z}b \\ | \\ a \end{array}} \end{array} ,$$

$$R^2f = \begin{array}{c} \boxed{\begin{array}{c} b \\ | \\ w(\bar{x}+\bar{y})b \\ / \quad \backslash \\ a \cdot \overline{x+y} \cdot x(\bar{x}+\bar{y})b \quad a \cdot \overline{x+y} \cdot y(\bar{x}+\bar{y})b \\ \backslash \quad / \\ a\bar{b}z \cdot \overline{x+y} \cdot b \\ | \\ a \end{array}} \end{array} ,$$

$$R^3f = \begin{array}{c} \boxed{\begin{array}{c} b \\ | \\ a\bar{w}b \\ / \quad \backslash \\ \dots \quad a\bar{b}z \cdot \overline{x+y} \cdot \overline{\bar{x}+\bar{y}} \cdot \bar{y} \cdot (x+y)\bar{w}b \\ \backslash \quad / \\ a\bar{b} \cdot \overline{\bar{x}+\bar{y}} \cdot \bar{w}b \\ | \\ a \end{array}} \end{array} ,$$

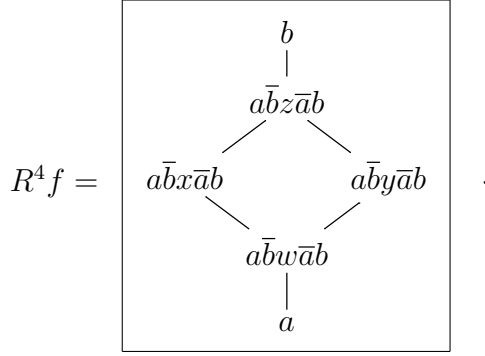
$$R^4 f = \begin{array}{c} \begin{array}{c} b \\ | \\ \bar{a}\bar{b}z\bar{a}\bar{b} \\ / \quad \backslash \\ \cdots \quad \bar{a}\bar{b} \cdot \overline{x+y} \cdot \overline{x+y} \cdot y \cdot (\overline{x+y}) \cdot (x+y) \cdot \bar{a}\bar{b} \\ \backslash \quad / \\ \bar{a}\bar{b}w\bar{a}\bar{b} \\ | \\ a \end{array} \end{array} .$$

Here, we have omitted the label at $(2, 1)$ for both $R^3 f$ and $R^4 f$, since it can be obtained from the respective label at $(1, 2)$ by interchanging x with y (thanks to an obvious symmetry between $(1, 2)$ and $(2, 1)$).

The above might suggest that the labels get progressively more complicated as we apply R over and over. For a general poset P , this is indeed the case. However, for our poset $P = \{1, 2\} \times \{1, 2\}$, a surprising periodicity-like pattern emerges. Indeed, our above expressions for $R^2 f$, $R^3 f$, $R^4 f$ can be simplified as follows:

$$R^2 f = \begin{array}{c} \begin{array}{c} b \\ | \\ w(\overline{x+y})b \\ / \quad \backslash \\ \bar{a}\bar{y}b \quad \bar{a}\bar{x}b \\ \backslash \quad / \\ \bar{a}\bar{b}z \cdot \overline{x+y} \cdot b \\ | \\ a \end{array} \end{array} ,$$

$$R^3 f = \begin{array}{c} \begin{array}{c} b \\ | \\ \bar{a}\bar{w}b \\ / \quad \backslash \\ \bar{a}\bar{b}z \cdot \overline{x+y} \cdot \bar{y}\bar{w}b \quad \bar{a}\bar{b}z \cdot \overline{x+y} \cdot \bar{x}\bar{w}b \\ \backslash \quad / \\ \bar{a}\bar{b} \cdot \overline{x+y} \cdot \bar{w}b \\ | \\ a \end{array} \end{array} ,$$



Thus, the labels of $R^4 f$ are closely related to those of f : For each $v \in P$, we have

$$(R^4 f)(v) = \bar{a}\bar{b} \cdot f(v) \cdot \bar{a}\bar{b}.$$

(This holds for $v = 0$ and $v = 1$ as well, as one can easily check.) Note that if $ab = ba$, then this entails that $(R^4 f)(v)$ is conjugate to v in \mathbb{K} .

In Theorem 4.7, we will generalize this phenomenon to arbitrary “rectangular” posets – i.e., posets of the form $\{1, 2, \dots, p\} \times \{1, 2, \dots, q\}$ with entrywise order. The “period” in this situation will be $p + q$.

Our $P = \{1, 2\} \times \{1, 2\}$ example also exhibits a reciprocity-like phenomenon. Indeed, our above expressions for Rf , $R^2 f$, $R^3 f$ reveal that

$$\begin{aligned} (Rf)(1, 1) &= a\bar{z}b = a \cdot \overline{f(2, 2)} \cdot b; \\ (R^2 f)(1, 2) &= a\bar{x}b = a \cdot \overline{f(2, 1)} \cdot b; \\ (R^2 f)(2, 1) &= a\bar{y}b = a \cdot \overline{f(1, 2)} \cdot b; \\ (R^3 f)(2, 2) &= a\bar{w}b = a \cdot \overline{f(1, 1)} \cdot b. \end{aligned}$$

These equalities relate the label of $R^{i+j-1} f$ at an element (i, j) with the label of f at the element $(3 - i, 3 - j)$ (which is, visually speaking, the “antipode” of the former element (i, j) on the Hasse diagram of P). To be specific, they say that

$$(R^{i+j-1} f)(i, j) = a \cdot \overline{f(3 - i, 3 - j)} \cdot b$$

for any $(i, j) \in P$. This too can be generalized to arbitrary rectangles (Theorem 4.8).

In the above calculation, we used the linear extension $((1, 1), (1, 2), (2, 1), (2, 2))$ of P to compute R as $T_{(1,1)} \circ T_{(1,2)} \circ T_{(2,1)} \circ T_{(2,2)}$. We could have just as well used the linear extension $((1, 1), (2, 1), (1, 2), (2, 2))$, obtaining the same result. But we could not have used the list $((1, 1), (1, 2), (2, 2), (2, 1))$ (for example), since it is not a linear extension (and indeed, $T_{(1,1)} \circ T_{(1,2)} \circ T_{(2,2)} \circ T_{(2,1)}$ would not give rise to any similar phenomenon).

This example shows that birational rowmotion behaves unexpectedly well for some posets. There are also some more serious motivations to study it: Birational rowmotion for

commutative \mathbb{K} generalizes Schützenberger’s classical “promotion” map on semistandard tableaux (see [GriRob14, Remark 11.6]), and is closely related to the Zamolodchikov periodicity conjecture in type AA (see [Roby15, §4.4]). The case of a noncommutative ring \mathbb{K} appears more baroque, but we expect it to find a combinatorial meaning sooner or later.

Before we formalize and prove the above phenomena, we first consider some general properties of R . We begin with an implicit description of birational rowmotion that does not involve toggles (but is essentially a restatement of Definition 3.16):

Proposition 3.18. Let $v \in P$. Let $f \in \mathbb{K}^{\widehat{P}}$. Assume that $Rf \neq \perp$. Then

$$(Rf)(v) = \left(\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u) \right) \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > v}} (Rf)(u)}.$$

Proof. This is merely the noncommutative analogue of [GriRob16, Proposition 19], and the proof in [GriRob16] can be used with straightforward modifications. \square

The following near-trivial fact completes the picture:

Proposition 3.19. Let $f \in \mathbb{K}^{\widehat{P}}$. Assume that $Rf \neq \perp$. Then $(Rf)(0) = f(0)$ and $(Rf)(1) = f(1)$.

Proof. None of the toggles T_v , when applied to a \mathbb{K} -labelling, changes the label of 0 or the label of 1. Hence, the same is true for the partial map R (since R is a composition of such toggles T_v). \square

A trivial corollary of Proposition 3.19 is:

Corollary 3.20. Let $f \in \mathbb{K}^{\widehat{P}}$ and $\ell \in \mathbb{N}$. Assume that $R^\ell f \neq \perp$. Then $(R^\ell f)(0) = f(0)$ and $(R^\ell f)(1) = f(1)$.

(Recall that \mathbb{N} denotes the set $\{0, 1, 2, \dots\}$ in this paper.)

3.6. Well-definedness lemmas

We next show some simple lemmas which say that certain inverses exist under the assumption that $R^\ell f$ is well-defined for some values of ℓ . These lemmas are easy and unexciting, but are necessary in order to rigorously prove the more substantial results that will follow. We recommend the reader skip the proofs, at least on a first reading.

▮ **Lemma 3.21.** Let $f \in \mathbb{K}^{\widehat{P}}$ and $k, \ell \in \mathbb{N}$ satisfy $k \leq \ell$ and $R^\ell f \neq \perp$. Then, $R^k f \neq \perp$.

Proof. We have $R^{\ell-k}(R^k f) = R^\ell f \neq \perp = R^{\ell-k}(\perp)$, so that $R^k f \neq \perp$. \square

▮ **Lemma 3.22.** Let $f \in \mathbb{K}^{\widehat{P}}$ satisfy $Rf \neq \perp$. Let $v \in P$. Then, $f(v)$ is invertible.

Proof. Proposition 3.18 yields

$$(Rf)(v) = \left(\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u) \right) \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > v}} (Rf)(u)}.$$

Thus, in particular, $\overline{f(v)}$ is well-defined. In other words, $f(v)$ is invertible. This proves Lemma 3.22. \square

▮ **Lemma 3.23.** Assume that $P \neq \emptyset$. Let $f \in \mathbb{K}^{\widehat{P}}$ satisfy $Rf \neq \perp$. Then, $f(1)$ is invertible.

Proof. We have $P \neq \emptyset$. Thus, the poset P has a maximal element y (by Proposition 1.9 (b)). This y then satisfies $1 \succ y$ in \widehat{P} .

We have $Rf \neq \perp$. Therefore, Proposition 3.18 (applied to $v = y$) yields

$$(Rf)(y) = \left(\sum_{\substack{u \in \widehat{P}; \\ u < y}} f(u) \right) \cdot \overline{f(y)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > y}} (Rf)(u)}.$$

Hence, in particular, $\overline{(Rf)(u)}$ is well-defined for each $u \in \widehat{P}$ satisfying $u \succ y$. We can apply this to $u = 1$ (since $1 \succ y$), and thus conclude that $\overline{(Rf)(1)}$ is well-defined. In other words, $(Rf)(1)$ is invertible. However, Proposition 3.19 yields $(Rf)(1) = f(1)$. Thus, $f(1)$ is invertible. \square

▮ **Lemma 3.24.** Assume that $P \neq \emptyset$. Let $f \in \mathbb{K}^{\widehat{P}}$ satisfy $R^2 f \neq \perp$. Then, $f(0)$ and $f(1)$ are invertible.

Proof. The poset P has a minimal element x (by Proposition 1.9 (a)).

From $R^2 f \neq \perp$, we obtain $Rf \neq \perp$ (by Lemma 3.21); thus, $Rf \in \mathbb{K}^{\widehat{P}}$. Hence, Lemma 3.23 yields that $f(1)$ is invertible. Furthermore, Lemma 3.22 (applied to Rf and x instead of f and v) yields that $(Rf)(x)$ is invertible (since $R(Rf) \neq \perp$).

Recall again that $Rf \neq \perp$. Hence, Proposition 3.18 (applied to $v = x$) yields

$$(Rf)(x) = \left(\sum_{\substack{u \in \widehat{P}; \\ u < x}} f(u) \right) \cdot \overline{f(x)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u > x}} (Rf)(u)}. \quad (2)$$

The only $u \in \widehat{P}$ satisfying $u \triangleleft x$ is the element 0 of \widehat{P} (since x is a minimal element of P). Thus, $\sum_{\substack{u \in \widehat{P}; \\ u \triangleleft x}} f(u) = f(0)$. Hence, (2) rewrites as

$$(Rf)(x) = f(0) \cdot \overline{f(x)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u \triangleright x}} (Rf)(u)}.$$

Solving this equality for $f(0)$, we obtain

$$f(0) = (Rf)(x) \cdot \left(\sum_{\substack{u \in \widehat{P}; \\ u \triangleright x}} (Rf)(u) \right) \cdot f(x).$$

The right hand side of this equality is a product of three invertible elements (indeed, the two factors $\sum_{\substack{u \in \widehat{P}; \\ u \triangleright x}} (Rf)(u)$ and $f(x)$ are invertible because their inverses appear in (2), and we already know that the factor $(Rf)(x)$ is invertible), and thus itself invertible. Hence, the left hand side is invertible. In other words, $f(0)$ is invertible. \square

Lemma 3.25. Let $f \in \mathbb{K}^{\widehat{P}}$ satisfy $Rf \neq \perp$. Let $v \in P$. Assume that v is not a minimal element of P . Then, $(Rf)(v)$ is invertible.

Proof. The element v of P is not minimal. Thus, there exists at least one element $w \in P$ satisfying $v \triangleright w$ ¹¹. Consider this w . Proposition 3.18 (applied to w instead of v) yields

$$(Rf)(w) = \left(\sum_{\substack{u \in \widehat{P}; \\ u \triangleleft w}} f(u) \right) \cdot \overline{f(w)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u \triangleright w}} (Rf)(u)}.$$

In particular, $\overline{(Rf)(u)}$ is well-defined for each $u \in \widehat{P}$ satisfying $u \triangleright w$. Applying this to $u = v$, we conclude that $\overline{(Rf)(v)}$ is well-defined (since $v \in P \subseteq \widehat{P}$ and $v \triangleright w$). In other words, $(Rf)(v)$ is invertible. \square

¹¹*Proof.* The element v of P is not minimal. Thus, there exists some $u \in P$ satisfying $u < v$. In other words, the set $P_{<v} := \{u \in P \mid u < v\}$ is nonempty. Consider this set $P_{<v}$ as a subposet of P (with its partial order inherited from P). Then, $P_{<v}$ is a finite poset and thus has a maximal element (by Proposition 1.9 (b), applied to $P_{<v}$ instead of P). Let m be this maximal element. Then, $m \in P_{<v} = \{u \in P \mid u < v\}$; in other words, $m \in P$ and $m < v$. There exists no element of $P_{<v}$ that is larger than m (since m is a **maximal** element of $P_{<v}$).

If there was some $w \in P$ satisfying $m < w < v$, then this w would belong to $P_{<v}$ (since $w < v$) but would be larger than m in the poset $P_{<v}$ (since $m < w$); this would contradict the fact that there exists no element of $P_{<v}$ that is larger than m . Hence, there is no $w \in P$ satisfying $m < w < v$. In other words, we have $m \triangleleft v$ in P (since $m < v$). In other words, we have $v \triangleright m$ in P . Hence, there exists at least one element $w \in P$ satisfying $v \triangleright w$ (namely, $w = m$).

Lemma 3.26. Assume that $P \neq \emptyset$. Let $f \in \mathbb{K}^{\widehat{P}}$ satisfy $Rf \neq \perp$. Let $v \in \widehat{P}$. Assume that $f(0)$ is invertible. Then, $(Rf)(v)$ is invertible.

Proof. If $v = 0$, then the claim follows from our assumption about $f(0)$ (since Proposition 3.19 yields $(Rf)(0) = f(0)$). If $v = 1$, then it instead follows from Lemma 3.23 (since Proposition 3.19 yields $(Rf)(1) = f(1)$). Thus, we assume from now on that v is neither 0 nor 1. Hence, $v \in P$.

If v is not a minimal element of P , then the claim follows from Lemma 3.25. Hence, we assume from now on that v is a minimal element of P . Therefore, the only $u \in \widehat{P}$ satisfying $u \triangleleft v$ is the element 0. Thus, $\sum_{\substack{u \in \widehat{P}; \\ u \triangleleft v}} f(u) = f(0)$. Now, Proposition 3.18 yields

$$(Rf)(v) = \underbrace{\left(\sum_{\substack{u \in \widehat{P}; \\ u \triangleleft v}} f(u) \right)}_{=f(0)} \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u \triangleright v}} (Rf)(u)} = f(0) \cdot \overline{f(v)} \cdot \overline{\sum_{\substack{u \in \widehat{P}; \\ u \triangleright v}} (Rf)(u)}.$$

The right hand side of this equality is a product of three invertible elements (since $f(0)$ is invertible, and since $\overline{f(v)}$ and $\overline{\sum_{\substack{u \in \widehat{P}; \\ u \triangleright v}} (Rf)(u)}$ are invertible¹²), and thus itself is invertible.

Thus, the left hand side is invertible as well. In other words, $(Rf)(v)$ is invertible. \square

4. The rectangle: statements of the results

4.1. The $p \times q$ -rectangle

As promised, we now state the phenomena observed in Example 3.17 in greater generality (and afterwards prove them). First we define the posets on which these phenomena manifest:

Definition 4.1. For $p \in \mathbb{Z}$, we let $[p]$ denote the totally ordered set $\{1, 2, \dots, p\}$ (with its usual total order: $1 < 2 < \dots < p$). This set is empty if $p \leq 0$.

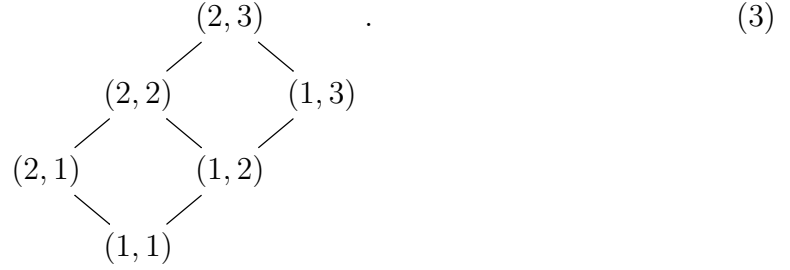
Definition 4.2. Let p and q be two positive integers. The $p \times q$ -rectangle will mean the Cartesian product $[p] \times [q]$ of the two posets $[p]$ and $[q]$. Explicitly, this is the set $[p] \times [q] = \{1, 2, \dots, p\} \times \{1, 2, \dots, q\}$, equipped with the partial order defined as follows: For two elements (i, j) and (i', j') of $[p] \times [q]$, we set $(i, j) \leq (i', j')$ if and only if $(i \leq i'$ and $j \leq j')$.

Henceforth, if we speak of $[p] \times [q]$, we implicitly assume that p and q are two positive integers.

The $p \times q$ -rectangle has been denoted by $\text{Rect}(p, q)$ in [GriRob14].

¹²because an inverse is always invertible

Example 4.3. Here is the Hasse diagram of the 2×3 -rectangle $[2] \times [3]$:



Convention 4.4. In the following, the Hasse diagram of a $p \times q$ -rectangle will always be drawn as in (3). That is, the elements (i, j) of $[p] \times [q]$ will be aligned in a rectangular grid, with the x -axis going southeast to northwest and the y -axis going southwest to northeast. Thus, for instance, the northwestern neighbor of an element (i, j) is always $(i + 1, j)$.

Two elements s and t of \widehat{P} will be called *adjacent* if they satisfy $s \succ t$ or $t \succ s$.

The poset $[p] \times [q]$ has a unique minimal element, $(1, 1)$, and a unique maximal element, (p, q) . Its covering relation can be characterized by the following easy remark (which will be used without explicit mention):

Remark 4.5. Let (i, j) and (i', j') be two elements of $[p] \times [q]$. Then, $(i, j) \prec (i', j')$ if and only if (i', j') is either $(i + 1, j)$ or $(i, j + 1)$.

Convention 4.6. Let $P = [p] \times [q]$. If f is a function defined on P or on \widehat{P} , and if (i, j) is any element of P , then we will write $f(i, j)$ for $f((i, j))$.

4.2. Periodicity

The following theorem (conjectured by the first author in 2014) generalizes the periodicity-like phenomenon seen in Example 3.17:

Theorem 4.7 (Periodicity theorem for the $p \times q$ -rectangle). Let $P = [p] \times [q]$, and let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^{p+q}f \neq \perp$. Set $a = f(0)$ and $b = f(1)$. Then, a and b are invertible, and for any $x \in \widehat{P}$ we have

$$(R^{p+q}f)(x) = a\bar{b} \cdot f(x) \cdot \bar{a}b. \tag{4}$$

If the ring \mathbb{K} is commutative¹³, then (4) simplifies to $(R^{p+q}f)(x) = f(x)$ (since commutativity of \mathbb{K} yields $a\bar{b} \cdot f(x) \cdot \bar{a}b = \underbrace{a\bar{a}}_{=1} \cdot f(x) \cdot \underbrace{\bar{b}b}_{=1} = f(x)$). Thus, if \mathbb{K} is commutative,

¹³or, more generally, if the elements a and b commute with each other and with $f(x)$

then the claim of Theorem 4.7 can be rewritten as $R^{p+q}f = f$, generalizing the main part of [GriRob14, Theorem 11.5] (which itself generalizes similar properties of rowmotion operators on other levels). Unlike in [GriRob14, Theorem 11.5], we cannot honestly claim that $R^{p+q} = \text{id}$ even when \mathbb{K} is commutative, since the partial map R^{p+q} takes the value \perp on some \mathbb{K} -labellings f (while id does not).

The parallel result for birational antichain rowmotion [JosRob20, Conjecture 5.10] follows from Theorem 4.7.

4.3. Reciprocity

Theorem 4.7 shows that the “periodicity phenomenon” we have observed on $[2] \times [2]$ in Example 3.17 was not a coincidence. The “reciprocity phenomenon” is similarly the $p = q = 2$ case of a general fact:

Theorem 4.8 (Reciprocity theorem for $p \times q$ -rectangle). Let $P = [p] \times [q]$. Fix $\ell \in \mathbb{N}$, and let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^\ell f \neq \perp$. Set $a = f(0)$ and $b = f(1)$. Let $(i, j) \in P$ satisfy $\ell - i - j + 1 \geq 0$. Then,

$$(R^\ell f)(i, j) = a \cdot \overline{(R^{\ell-i-j+1}f)(p+1-i, q+1-j)} \cdot b. \quad (5)$$

Theorem 4.8 directly generalizes the analogous theorem [GriRob14, Theorem 11.7] in the commutative setting.

5. Twisted periodicity follows from reciprocity

Our first step towards the proofs of twisted periodicity (Theorem 4.7) and reciprocity (Theorem 4.8) is to show that the latter implies the former.

Proof of Theorem 4.7 using Theorem 4.8. Assume that Theorem 4.8 has been proved. Let p, q, P, f, a and b be as in Theorem 4.7. Let $x \in \widehat{P}$. From $p \geq 1$ and $q \geq 1$, we obtain $p + q \geq 2$. Hence, from $R^{p+q}f \neq \perp$, we obtain $R^2f \neq \perp$ (by Lemma 3.21). Therefore, Lemma 3.24 yields that a and b are invertible (since $a = f(0)$ and $b = f(1)$).

It remains to prove (4). First, we note that

$$a\bar{b} \cdot \underbrace{f(0)}_{=a} \cdot \bar{a}b = a\bar{b} \cdot \underbrace{a \cdot \bar{a}}_{=1} b = a \underbrace{\bar{b} \cdot b}_{=1} = a.$$

However, Corollary 3.20 yields $(R^{p+q}f)(0) = f(0) = a$. Comparing these, we find that $(R^{p+q}f)(0) = a\bar{b} \cdot f(0) \cdot \bar{a}b$. Thus, the equality (4) holds for $x = 0$. Similarly, this equality also holds for $x = 1$. So from now on, we WLOG assume that x is neither 0 nor 1. Hence, $x = (i, j) \in P = [p] \times [q]$.

From $R^{p+q}f \neq \perp$, we obtain $Rf \neq \perp$ (by Lemma 3.21, since $1 \leq 2 \leq p + q$). Thus, Lemma 3.22 (applied to $v = (i, j)$) yields that $f(i, j)$ is invertible. Hence, $\overline{f(i, j)}$ is well-defined. The element $\overline{f(i, j)}$ of \mathbb{K} is invertible (since it has inverse $f(i, j)$).

Set $i' := p + 1 - i \in [p]$ and $j' := q + 1 - j \in [q]$, so that $(i', j') \in [p] \times [q] = P$ and $i' + j' \geq 1 + 1 = 2$. Also, the definitions of i' and j' readily yield $p + 1 - i' = i$ and $q + 1 - j' = j$ and $i' + j' - 1 = p + q - i - j + 1$.

Now $i' + j' \leq p + q$, so $i' + j' - 1 \leq i' + j' \leq p + q$ and thus $R^{i'+j'-1}f \neq \perp$ (by Lemma 3.21, since $R^{p+q}f \neq \perp$). Thus, Theorem 4.8 (applied to $i' + j' - 1$, i' and j' instead of ℓ , i and j) yields

$$\begin{aligned} (R^{i'+j'-1}f)(i', j') &= a \cdot \overbrace{\left(R^{(i'+j'-1)-i'-j'+1}f \right)}_{\substack{=f \\ \text{(since } (i'+j'-1)-i'-j'+1=0)}} \left(\underbrace{p+1-i'}_{=i}, \underbrace{q+1-j'}_{=j} \right) \cdot b \\ &= a \cdot \overline{f(i, j)} \cdot b. \end{aligned} \tag{6}$$

However, we also have $p + q - i - j + 1 = i' + j' - 1 \geq 0$ (since $i' + j' \geq 2 \geq 1$). Thus, Theorem 4.8 (applied to $\ell = p + q$) yields

$$\begin{aligned} (R^{p+q}f)(i, j) &= a \cdot \overline{(R^{p+q-i-j+1}f)(p+1-i, q+1-j)} \cdot b \\ &= a \cdot \overline{(R^{i'+j'-1}f)(i', j')} \cdot b \quad \left(\begin{array}{l} \text{since } p+q-i-j+1 = i'+j'-1 \\ \text{and } p+1-i = i' \text{ and } q+1-j = j' \end{array} \right) \\ &= a \cdot \underbrace{\overline{a \cdot f(i, j) \cdot b}}_{\substack{= \bar{b} \cdot f(i, j) \cdot \bar{a} \\ \text{(by Proposition 2.3 (c),} \\ \text{since } a \text{ and } \overline{f(i, j)} \text{ and } b \text{ are invertible)}}} \cdot b \quad (\text{by (6)}) \\ &= \bar{a} \cdot f(i, j) \cdot \bar{a} b. \end{aligned}$$

Since $x = (i, j)$, we can rewrite this as

$$(R^{p+q}f)(x) = \bar{a} \cdot f(x) \cdot \bar{a} b.$$

Thus twisted periodicity (Theorem 4.7) is proved, assuming reciprocity (Theorem 4.8) holds. \square

6. Proof of reciprocity: notations

It now suffices to prove Theorem 4.8, which will be the ultimate goal of the next few sections. First we introduce some notations that will be used throughout these sections.

Fix two positive integers p and q . Assume that $P = [p] \times [q]$. Let $f \in \mathbb{K}^{\hat{P}}$ be a \mathbb{K} -labelling of P . Set

$$a := f(0) \quad \text{and} \quad b := f(1).$$

For any $x = (i, j) \in P$, we define an element $x^\sim \in P$ by

$$x^\sim := (p + 1 - i, q + 1 - j);$$

. we call this element x^\sim the *antipode* of x . Thus, the desired equality (5) can be rewritten as

$$(R^\ell f)(x) = a \cdot \overline{(R^{\ell-i-j+1}f)(x^\sim)} \cdot b \quad (7)$$

for $x = (i, j)$.

For any $x \in \widehat{P}$ and $\ell \in \mathbb{N}$, we write

$$x_\ell := (R^\ell f)(x), \quad (8)$$

which is well-defined whenever $R^\ell f \neq \perp$. This compact notation will make upcoming formulas more readable.

In particular, for each $x \in \widehat{P}$, we have $x_0 = (R^0 f)(x) = f(x)$. Moreover, for each $\ell \in \mathbb{N}$ satisfying $R^\ell f \neq \perp$, we have

$$0_\ell = (R^\ell f)(0) = a \quad (\text{via Corollary 3.20}) \quad (9)$$

and similarly $1_\ell = b$.

We can further rewrite the equality (7) as $x_\ell = a \cdot \overline{x_{\ell-i-j+1}^\sim} \cdot b$ (since $x_\ell = (R^\ell f)(x)$ and $x_{\ell-i-j+1}^\sim = (R^{\ell-i-j+1}f)(x^\sim)$). Hence, our desired Theorem 4.8 takes the following form:

Theorem 4.8, restated. If $x = (i, j) \in P$ and $\ell \in \mathbb{N}$ satisfy $\ell - i - j + 1 \geq 0$ and $R^\ell f \neq \perp$, then

$$x_\ell = a \cdot \overline{x_{\ell-i-j+1}^\sim} \cdot b. \quad (10)$$

Proposition 3.18 yields that for each $v \in P$, we have¹⁴

$$(Rf)(v) = \left(\sum_{u < v} f(u) \right) \cdot \overline{f(v)} \cdot \overline{\sum_{u > v} (Rf)(u)}. \quad (11)$$

(In both sums, u ranges over \widehat{P} ; from now on, this will always be understood if not otherwise specified.) Applying this equality (11) to $R^\ell f$ instead of f , we obtain

$$(R^{\ell+1}f)(v) = \left(\sum_{u < v} (R^\ell f)(u) \right) \cdot \overline{(R^\ell f)(v)} \cdot \overline{\sum_{u > v} (R^{\ell+1}f)(u)}$$

for each $v \in P$ and $\ell \in \mathbb{N}$ satisfying $R^{\ell+1}f \neq \perp$ (since $R(R^\ell f) = R^{\ell+1}f$).

Using (8), we can rewrite this as follows:

$$v_{\ell+1} = \left(\sum_{u < v} u_\ell \right) \cdot \overline{v_\ell} \cdot \overline{\sum_{u > v} u_{\ell+1}} \quad (12)$$

for each $v \in P$ and $\ell \in \mathbb{N}$ satisfying $R^{\ell+1}f \neq \perp$.

¹⁴assuming that $Rf \neq \perp$

Next, we formally define the paths that will play a key role in the proof. A *path* means a sequence (v_0, v_1, \dots, v_k) of elements of \widehat{P} satisfying $v_0 \succ v_1 \succ \dots \succ v_k$. We denote this path by $(v_0 \succ v_1 \succ \dots \succ v_k)$, and we will call it a *path from v_0 to v_k* (or, for short, a *path $v_0 \rightarrow v_k$*). The *vertices* of this path are defined to be the elements v_0, v_1, \dots, v_k . We say that this path *starts* at v_0 and *ends* at v_k .

For each $v \in P$ and $\ell \in \mathbb{N}$, we set¹⁵

$$A_\ell^v := v_\ell \cdot \overline{\sum_{u < v} u_\ell} \quad \text{and} \quad V_\ell^v := \overline{\sum_{u > v} u_\ell} \cdot \overline{v_\ell}.$$

¹⁶ Furthermore, when $v \in \{0, 1\}$, we set

$$A_\ell^v := 1 \quad \text{and} \quad V_\ell^v := 1 \quad (13)$$

for all $\ell \in \mathbb{N}$.

For any path $\mathbf{p} = (v_0 \succ v_1 \succ \dots \succ v_k)$ and any $\ell \in \mathbb{N}$, we set

$$\begin{aligned} A_\ell^{\mathbf{p}} &:= A_\ell^{v_0} A_\ell^{v_1} \dots A_\ell^{v_k} & \text{and} \\ V_\ell^{\mathbf{p}} &:= V_\ell^{v_0} V_\ell^{v_1} \dots V_\ell^{v_k} \end{aligned}$$

(assuming that the factors on the right hand sides are well-defined).

If u and v are elements of \widehat{P} , and if $\ell \in \mathbb{N}$, then we set

$$A_\ell^{u \rightarrow v} := \sum_{\mathbf{p} \text{ is a path from } u \text{ to } v} A_\ell^{\mathbf{p}} \quad \text{and} \quad (14)$$

$$V_\ell^{u \rightarrow v} := \sum_{\mathbf{p} \text{ is a path from } u \text{ to } v} V_\ell^{\mathbf{p}} \quad (15)$$

(assuming that all addends on the right hand sides are well-defined).

Example 6.1. Let $P = [2] \times [2]$ and $f \in \mathbb{K}^{\widehat{P}}$ be as in Example 3.17. Then,

$$\begin{aligned} (1, 1)^\sim &= (2, 2), & (1, 2)^\sim &= (2, 1), & (2, 1)^\sim &= (1, 2), & (2, 2)^\sim &= (1, 1), \\ (1, 1)_0 &= f(1, 1) = w, & (1, 1)_1 &= (Rf)(1, 1) = a\bar{z}b, \\ (1, 1)_2 &= (R^2f)(1, 1) = a\bar{b}z \cdot \overline{x + y} \cdot b, \\ (1, 2)_2 &= (R^2f)(1, 2) = a \cdot \overline{x + y} \cdot y \cdot (\overline{x + y}) b. \end{aligned}$$

There are only two paths from $(2, 2)$ to $(1, 1)$: namely, $((2, 2) \succ (1, 2) \succ (1, 1))$ and $((2, 2) \succ (2, 1) \succ (1, 1))$. Each of these two paths has three vertices. There are no paths

¹⁵We recall that the summation signs “ $\sum_{u < v}$ ” and “ $\sum_{u > v}$ ” mean “ $\sum_{\substack{u \in \widehat{P}; \\ u < v}}$ ” and “ $\sum_{\substack{u \in \widehat{P}; \\ u > v}}$ ”, respectively.

¹⁶These elements A_ℓ^v and V_ℓ^v are not always well-defined. For A_ℓ^v to be well-defined, we need to have $R^\ell f \neq \perp$, and we need the element $\sum_{u < v} u_\ell$ to be invertible. For V_ℓ^v to be well-defined, we need to have $R^\ell f \neq \perp$, and we need the elements $\overline{u_\ell}$ (for $u > v$) and $\sum_{u > v} \overline{u_\ell}$ and v_ℓ to be invertible.

from $(1, 1)$ to $(2, 2)$, since we don't have $(1, 1) \geq (2, 2)$. The only path from 0 to 0 is the trivial path (0) .

We have

$$\begin{aligned}
A_0^{(1,1)} &= (1, 1)_0 \cdot \overline{\sum_{u \prec (1,1)} u_0} = (1, 1)_0 \cdot \overline{0_0} = w \cdot \bar{a}, \\
A_0^{(2,2)} &= (2, 2)_0 \cdot \overline{\sum_{u \prec (2,2)} u_0} = (2, 2)_0 \cdot \overline{(1, 2)_0 + (2, 1)_0} = z \cdot \overline{y + x}, \\
A_1^{(1,1)} &= (1, 1)_1 \cdot \overline{\sum_{u \prec (1,1)} u_1} = (1, 1)_1 \cdot \overline{0_1} = a\bar{z}b \cdot \bar{a}, \\
V_0^{(1,1)} &= \overline{\sum_{u \succ (1,1)} \bar{u}_0} \cdot \overline{(1, 1)_0} = \overline{(1, 2)_0 + (2, 1)_0} \cdot \overline{(1, 1)_0} = \overline{y + x} \cdot \bar{w}, \\
V_1^{(1,1)} &= \overline{\sum_{u \succ (1,1)} \bar{u}_1} \cdot \overline{(1, 1)_1} = \overline{(1, 2)_1 + (2, 1)_1} \cdot \overline{(1, 1)_1} \\
&= \overline{w\bar{y}(x+y)\bar{z}b + w\bar{x}(x+y)\bar{z}b} \cdot \overline{a\bar{z}b} = w \cdot \bar{a} \quad (\text{after simplifications}).
\end{aligned}$$

Furthermore, for any $\ell \in \mathbb{N}$, we have

$$\begin{aligned}
A_\ell^{((2,2) \succ (1,2) \succ (1,1))} &= A_\ell^{(2,2)} A_\ell^{(1,2)} A_\ell^{(1,1)}; \\
A_\ell^{(2,2) \rightarrow (1,1)} &= A_\ell^{((2,2) \succ (1,2) \succ (1,1))} + A_\ell^{((2,2) \succ (2,1) \succ (1,1))} \\
&= A_\ell^{(2,2)} A_\ell^{(1,2)} A_\ell^{(1,1)} + A_\ell^{(2,2)} A_\ell^{(2,1)} A_\ell^{(1,1)}
\end{aligned}$$

(and similarly for V instead of A).

The letter ℓ will always stand for a nonnegative integer (but will not be fixed).

7. Proof of reciprocity: simple lemmas

Throughout this section, we use the notations introduced in Section 6.

Let us prove some relations between the elements we have introduced. We begin with a well-definedness result:

Lemma 7.1. Let $\ell \in \mathbb{N}$ be such that $\ell \geq 1$ and $R^\ell f \neq \perp$. Assume furthermore that a is invertible. Let $v \in \widehat{P}$. Then:

- (a) The element v_ℓ is well-defined and invertible.
- (b) The element $v_{\ell-1}$ is well-defined and invertible.
- (c) The element $A_{\ell-1}^v$ is well-defined and invertible.
- (d) The element V_ℓ^v is well-defined and invertible.

Proof. From $R^\ell f \neq \perp$, we obtain $R^{\ell-1} f \neq \perp$. Hence, Corollary 3.20 yields that $(R^{\ell-1} f)(0) = f(0) = a$, which is invertible by assumption.

(a) Clearly, $v_\ell = (R^\ell f)(v)$ is well-defined, and we have $v_\ell = (R(R^{\ell-1} f))(v)$. Hence, Lemma 3.26 (applied to $R^{\ell-1} f$ instead of f) yields that v_ℓ is invertible.

(b) If $v = 0$, then this follows from part (a), because (9) yields that $v_{\ell-1} = a = v_\ell$ in this case. An analogous argument works if $v = 1$. Thus, we WLOG assume that $v \notin \{0, 1\}$, so that $v \in P$. The element $v_{\ell-1} = (R^{\ell-1} f)(v)$ is clearly well-defined, and is invertible by Lemma 3.22 (applied to $R^{\ell-1} f$ instead of f).

(c) If $v \in \{0, 1\}$, then this follows from (13). Otherwise, $v \in P$. Applying (12) to $\ell - 1$ instead of ℓ , we obtain

$$v_\ell = \left(\sum_{u < v} u_{\ell-1} \right) \cdot \overline{v_{\ell-1}} \cdot \overline{\sum_{u > v} \overline{u_\ell}}. \quad (16)$$

This equality shows that $\overline{v_{\ell-1}}$ and $\overline{\sum_{u > v} \overline{u_\ell}}$ are well-defined, i.e., the elements $v_{\ell-1}$ and $\sum_{u > v} \overline{u_\ell}$ are invertible. Also, v_ℓ is invertible (by Lemma 7.1 (a)).

Solving the equality (16) for the first factor on its right hand side, we obtain

$$\sum_{u < v} u_{\ell-1} = v_\ell \cdot \left(\sum_{u > v} \overline{u_\ell} \right) \cdot v_{\ell-1}.$$

The right hand side of this equality is a product of three invertible elements; thus, both sides are invertible. Therefore, the element $\overline{\sum_{u < v} u_{\ell-1}}$ is well-defined, hence invertible (since an inverse is always invertible).

Finally, $A_{\ell-1}^v$ is defined to be the product $v_{\ell-1} \cdot \overline{\sum_{u < v} u_{\ell-1}}$, and thus is well-defined and invertible because both of its factors are.

(d) If $v \in \{0, 1\}$, then this follows from (13). Otherwise, $v \in P$.

Lemma 7.1 (a) shows that v_ℓ is invertible; hence, $\overline{v_\ell}$ is well-defined, and invertible. Also, in the proof of Lemma 7.1 (c), we have shown that $\overline{\sum_{u > v} \overline{u_\ell}}$ is well-defined, so it too is invertible. Finally, V_ℓ^v is defined to be the product $\overline{\sum_{u > v} \overline{u_\ell}} \cdot \overline{v_\ell}$, and thus is also well-defined and invertible because both of its factors are. \square

Next we show some simple recursions for $A_\ell^{s \rightarrow t}$ and $V_\ell^{s \rightarrow t}$:

Proposition 7.2. Let s and t be two distinct elements of \widehat{P} , and fix $\ell \in \mathbb{N}$. Then

$$A_\ell^{s \rightarrow t} = A_\ell^s \sum_{\substack{u \in \widehat{P}; \\ s > u}} A_\ell^{u \rightarrow t} \quad (17)$$

$$= \sum_{\substack{u \in \widehat{P}; \\ u > t}} A_\ell^{s \rightarrow u} A_\ell^t \quad (18)$$

and

$$V_\ell^{s \rightarrow t} = V_\ell^s \sum_{\substack{u \in \widehat{P}; \\ s \succ u}} V_\ell^{u \rightarrow t} \quad (19)$$

$$= \sum_{\substack{u \in \widehat{P}; \\ u \succ t}} V_\ell^{s \rightarrow u} V_\ell^t. \quad (20)$$

Here, we assume that all the terms in the respective equalities are well-defined.

Proof. Since $s \neq t$, every path from s to t must contain an element covered by s as its second vertex.

Fix an element $u \in \widehat{P}$ satisfying $s \succ u$. If $(v_0 \succ v_1 \succ \cdots \succ v_k)$ is a path from s to t satisfying $v_1 = u$, then $(v_1 \succ v_2 \succ \cdots \succ v_k)$ is a path from u to t . Hence, we have found a map

$$\begin{aligned} & \text{from } \{\text{paths } (v_0 \succ v_1 \succ \cdots \succ v_k) \text{ from } s \text{ to } t \text{ satisfying } v_1 = u\} \\ & \text{to } \{\text{paths from } u \text{ to } t\} \end{aligned}$$

that sends each path $(v_0 \succ v_1 \succ \cdots \succ v_k)$ to $(v_1 \succ v_2 \succ \cdots \succ v_k)$. This map is a bijection (since any path from u to t can be uniquely extended to a path from s to t by inserting the vertex s at the front). We can use this bijection to substitute $(v_1 \succ v_2 \succ \cdots \succ v_k)$ for \mathbf{p} in a sum that ranges over all paths \mathbf{p} from u to t . In particular,

$$\begin{aligned} & \sum_{\mathbf{p} \text{ is a path from } u \text{ to } t} A_\ell^s A_\ell^{\mathbf{p}} \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_k) \text{ is a path from } s \text{ to } t; \\ v_1 = u}} \underbrace{A_\ell^s}_{=A_\ell^{v_0}} \underbrace{A_\ell^{(v_1 \succ v_2 \succ \cdots \succ v_k)}}_{=A_\ell^{v_1} A_\ell^{v_2} \cdots A_\ell^{v_k}} \\ & \quad \text{(since } s=v_0 \text{) (by the definition of } A_\ell^{(v_1 \succ v_2 \succ \cdots \succ v_k)}) \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_k) \text{ is a path from } s \text{ to } t; \\ v_1 = u}} \underbrace{A_\ell^{v_0} A_\ell^{v_1} A_\ell^{v_2} \cdots A_\ell^{v_k}}_{=A_\ell^{(v_0 \succ v_1 \succ \cdots \succ v_k)}} \\ & \quad \text{(by the definition of } A_\ell^{(v_0 \succ v_1 \succ \cdots \succ v_k)}) \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_k) \text{ is a path from } s \text{ to } t; \\ v_1 = u}} A_\ell^{(v_0 \succ v_1 \succ \cdots \succ v_k)}. \quad (21) \end{aligned}$$

Now, forget that we fixed u . We thus have proved (21) for each $u \in \widehat{P}$ satisfying $s \succ u$.

The definition of $A_\ell^{s \rightarrow t}$ yields

$$\begin{aligned}
A_\ell^{s \rightarrow t} &= \sum_{\substack{(v_0 \triangleright v_1 \triangleright \dots \triangleright v_k) \\ \text{is a path from } s \text{ to } t}} A_\ell^{(v_0 \triangleright v_1 \triangleright \dots \triangleright v_k)} \\
&= \sum_{\substack{u \in \widehat{P}; \\ s \triangleright u}} \sum_{\substack{(v_0 \triangleright v_1 \triangleright \dots \triangleright v_k) \\ \text{is a path from } s \text{ to } t; \\ v_1 = u}} A_\ell^{(v_0 \triangleright v_1 \triangleright \dots \triangleright v_k)} \\
&\quad \text{(because any path } (v_0 \triangleright v_1 \triangleright \dots \triangleright v_k) \text{ from } s \text{ to } t \text{ has a well-defined} \\
&\quad \text{second vertex } v_1, \text{ and this second vertex } v_1 \text{ satisfies } s \triangleright v_1) \\
&= \sum_{\substack{u \in \widehat{P}; \\ s \triangleright u}} \underbrace{\sum_{\substack{(v_0 \triangleright v_1 \triangleright \dots \triangleright v_k) \\ \text{is a path from } s \text{ to } t; \\ v_1 = u}} A_\ell^{(v_0 \triangleright v_1 \triangleright \dots \triangleright v_k)}}_{= \sum_{\substack{\mathbf{p} \text{ is a path from } u \text{ to } t \\ \text{(by (21))}}} A_\ell^s A_\ell^{\mathbf{p}}} \\
&= \sum_{\substack{u \in \widehat{P}; \\ s \triangleright u}} \sum_{\mathbf{p} \text{ is a path from } u \text{ to } t} A_\ell^s A_\ell^{\mathbf{p}} = A_\ell^s \sum_{\substack{u \in \widehat{P}; \\ s \triangleright u}} \underbrace{\sum_{\mathbf{p} \text{ is a path from } u \text{ to } t} A_\ell^{\mathbf{p}}}_{= A_\ell^{u \rightarrow t}} \\
&\quad \text{(by the definition of } A_\ell^{u \rightarrow t}) \\
&= A_\ell^s \sum_{\substack{u \in \widehat{P}; \\ s \triangleright u}} A_\ell^{u \rightarrow t}.
\end{aligned}$$

This proves (17). The same argument (but with each A symbol replaced by an V symbol) proves (19). Moreover, a similar argument (but now classifying paths from s to t according to their second-to-last vertex instead of their second vertex) establishes (18) and (20). Thus, Proposition 7.2 is proven. \square

The next proposition uses the products V_ℓ^v and $A_{\ell-1}^v$ to rewrite the equality (12) (which is essentially the definition of birational rowmotion) in a slick way:

Proposition 7.3 (Transition equation in $\mathbf{A-V}$ -form). Let $v \in \widehat{P}$ and $\ell \geq 1$ be such that $R^\ell f \neq \perp$. Assume that a is invertible. Then,

$$V_\ell^v = A_{\ell-1}^v.$$

Proof. If v is 0 or 1, then the equality $V_\ell^v = A_{\ell-1}^v$ holds because both of its sides are 1 (by (13)). Thus, we assume WLOG that $v \in P$.

Lemma 7.1 (a) yields that v_ℓ is well-defined and invertible, while Lemma 7.1 (c,d) yield that V_ℓ^v and $A_{\ell-1}^v$ are well-defined. Since $A_{\ell-1}^v$ is defined as $v_{\ell-1} \cdot \sum_{u < v} u_{\ell-1}$, this entails that

$\sum_{u < v} u_{\ell-1}$ is invertible.

If $\alpha, \beta, \gamma, \delta$ are four invertible elements of \mathbb{K} satisfying $\alpha = \beta\overline{\gamma\delta}$, then

$$\delta\overline{\alpha} = \delta \underbrace{\overline{\beta\overline{\gamma\delta}}}_{=\overline{\delta\gamma\beta}} = \underbrace{\overline{\delta\overline{\delta}}}_{=1} \overline{\gamma\beta} = \overline{\gamma\beta}. \quad (22)$$

Applying (12) to $\ell - 1$ instead of ℓ , we find

$$v_\ell = \left(\sum_{u < v} u_{\ell-1} \right) \cdot \overline{v_{\ell-1}} \cdot \overline{\sum_{u > v} \overline{u_\ell}}.$$

Thus, (22) (applied to $\alpha = v_\ell$, $\beta = \sum_{u < v} u_{\ell-1}$, $\gamma = v_{\ell-1}$ and $\delta = \overline{\sum_{u > v} \overline{u_\ell}}$) yields

$$\overline{\sum_{u > v} \overline{u_\ell}} \cdot \overline{v_\ell} = v_{\ell-1} \cdot \overline{\sum_{u < v} u_{\ell-1}}.$$

But the left hand side of this equality is V_ℓ^v (by the definition of V_ℓ^v), whereas the right hand side is $A_{\ell-1}^v$. Hence, this equality simplifies to $V_\ell^v = A_{\ell-1}^v$. This proves Proposition 7.3. \square

As a consequence of Proposition 7.3, we have:

Corollary 7.4. Let \mathbf{p} be a path. Let $\ell \geq 1$ be such that $R^\ell f \neq \perp$. Assume that a is invertible. Then,

$$V_\ell^{\mathbf{p}} = A_{\ell-1}^{\mathbf{p}}.$$

Corollary 7.5. Let $u, v \in \widehat{P}$. Let $\ell \in \mathbb{N}$ be such that $\ell \geq 1$ and $R^\ell f \neq \perp$. Assume that a is invertible. Then,

$$V_\ell^{u \rightarrow v} = A_{\ell-1}^{u \rightarrow v}. \quad (23)$$

The next theorem gives ways to recover the labels $u_\ell = (R^\ell f)(u)$ from some of the sums defined in (14) and (15).¹⁷

Theorem 7.6 (path formulas for rectangle). Let $\ell \in \mathbb{N}$. Assume that a is invertible. Then:

(a) If $R^\ell f \neq \perp$ and $\ell \geq 1$, then each $u \in P$ satisfies

$$u_\ell = \overline{V_\ell^{1 \rightarrow u}} \cdot b$$

(and the inverse $\overline{V_\ell^{1 \rightarrow u}}$ is well-defined).

(b) If $R^{\ell+1} f \neq \perp$, then each $u \in P$ satisfies

$$u_\ell = A_\ell^{u \rightarrow 0} \cdot a.$$

¹⁷The condition $\ell \geq 1$ in Theorem 7.6 (a) and (c) is meant to ensure that $V_\ell^{1 \rightarrow u}$ and $V_\ell^{(p,q) \rightarrow u}$ are invertible. It can be replaced by directly requiring the latter.

(c) If $R^\ell f \neq \perp$ and $\ell \geq 1$, then each $u \in P$ satisfies

$$u_\ell = \overline{V_\ell^{(p,q) \rightarrow u}} \cdot b$$

(and the inverse $\overline{V_\ell^{(p,q) \rightarrow u}}$ is well-defined).

(d) If $R^{\ell+1} f \neq \perp$, then each $u \in P$ satisfies

$$u_\ell = A_\ell^{u \rightarrow (1,1)} \cdot a.$$

Proof of Theorem 7.6. (a) Assume that $R^\ell f \neq \perp$ and $\ell \geq 1$. Then, Lemma 7.1 (d) yields that the element V_ℓ^v is well-defined and invertible for each $v \in \widehat{P}$. Hence, the element $V_\ell^{\mathbf{p}}$ is well-defined for each path \mathbf{p} . Therefore, the element $V_\ell^{1 \rightarrow u}$ is well-defined for each $u \in P$.

Next, we will prove the equality

$$V_\ell^{1 \rightarrow u} = b\overline{u_\ell} \quad \text{for each } u \in P. \quad (24)$$

(The $\overline{u_\ell}$ on the right hand side here is well-defined, since Lemma 7.1 (a) (applied to $v = u$) shows that u_ℓ is well-defined and invertible.)

Proof of (24). We utilize downwards induction on u . This is a version of strong induction in which we fix an element $v \in P$ and assume (as the induction hypothesis) that (24) holds for all $u \in P$ satisfying $u > v$. We will then prove that (24) also holds for $u = v$. Since the poset P is finite, this will entail that (24) holds for all $u \in P$.

Let $v \in P$. Assume (as the induction hypothesis) that (24) holds for all $u \in P$ satisfying $u > v$. In other words, we have $V_\ell^{1 \rightarrow u} = b\overline{u_\ell}$ for each $u \in P$ satisfying $u > v$. Thus, in particular, we have

$$V_\ell^{1 \rightarrow u} = b\overline{u_\ell} \quad \text{for each } u \in P \text{ satisfying } u \succ v. \quad (25)$$

Note also that the only path from 1 to 1 is the trivial path (1). Hence,

$$V_\ell^{1 \rightarrow 1} = V_\ell^{(1)} = V_\ell^1 = 1 = b\overline{1_\ell} \quad (26)$$

(since $1_\ell = b$).

However, $1 \neq v$ (since $1 \notin P$ and $v \in P$). Thus, (20) (applied to $s = 1$ and $t = v$)

yields

$$\begin{aligned}
V_\ell^{1 \rightarrow v} &= \sum_{\substack{u \in \hat{P}; \\ u > v}} V_\ell^{1 \rightarrow u} && V_\ell^v \\
&= \sum_{u > v} \underbrace{V_\ell^{1 \rightarrow u}}_{=b\bar{u}_\ell} && \text{(indeed, this follows from (25) when } u \in P, \\
&\quad \text{(since our sums range over } \hat{P} \text{ by default)} && \text{and follows from (26) when } u=1; \\
&\quad \text{and there are no other possibilities, since } u > v \text{ rules out } u=0) \\
&= \sum_{u > v} b\bar{u}_\ell V_\ell^v = b \left(\sum_{u > v} \bar{u}_\ell \right) \underbrace{V_\ell^v}_{= \sum_{u > v} \overline{\bar{u}_\ell \cdot \bar{v}_\ell}} && = b \left(\sum_{u > v} \bar{u}_\ell \right) \underbrace{\sum_{u > v} \overline{\bar{u}_\ell \cdot \bar{v}_\ell}}_{=1} = b\bar{v}_\ell. \\
&\quad \text{(by the definition of } V_\ell^v)
\end{aligned}$$

In other words, (24) holds for $u = v$. This completes the induction step. Thus, we have proved (24) by induction. \square

Note that 1_ℓ is invertible (by Lemma 7.1 (a), applied to $v = 1$). In other words, b is invertible (since $1_\ell = b$).

Now, let $u \in P$. Then, $b\bar{u}_\ell$ is invertible (since b and \bar{u}_ℓ are). In view of (24), this means that $V_\ell^{1 \rightarrow u}$ is invertible. Hence, $\overline{V_\ell^{1 \rightarrow u}}$ is well-defined. Solving (24) for u_ℓ , we thus obtain $u_\ell = \overline{V_\ell^{1 \rightarrow u}} \cdot b$. This proves Theorem 7.6 (a).

(b) This proof is rather similar to that of part (a), but uses upwards induction instead of downwards induction (and applies (17) instead of (20)).

(c) Let $u \in P$. Recall that (p, q) is the unique maximal element of P . Therefore, each path from 1 to u begins with the step $1 \succ (p, q)$. Thus, $V_\ell^{1 \rightarrow u} = V_\ell^{(p, q) \rightarrow u}$ (since $V_\ell^1 = 1$). Hence, part (c) follows from (a).

Similarly, part (d) follows from (b). \square

Remark 7.7. Proposition 7.2 and parts (a) and (b) of Theorem 7.6 hold more generally if P is replaced by any finite poset (not necessarily a rectangle). The proofs we gave above work in that generality. Parts (c) and (d) of Theorem 7.6 can be similarly generalized as long as the poset P has a global maximum (for part (c)) and a global minimum (for part (d)); all we need to do is to replace (p, q) by the global maximum and $(1, 1)$ by the global minimum. We will have no need for this generality, though.

8. Proof of reciprocity: the case $(i, j) = (1, 1)$

Now, we are mostly ready to prove that Theorem 4.8 holds in the case when $(i, j) = (1, 1)$. For reasons both technical and pedagogical, it is useful for us to dispose of this case now in order to have less work to do later. First, we prove Theorem 4.8 for $(i, j) = (1, 1)$ under the extra assumption that a is invertible:

Lemma 8.1. Assume that P is the $p \times q$ -rectangle $[p] \times [q]$. Let $\ell \in \mathbb{N}$ be such that $\ell \geq 1$. Let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^\ell f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Assume that a is invertible. Then,

$$(R^\ell f)(1, 1) = a \cdot \overline{(R^{\ell-1} f)(p, q)} \cdot b.$$

Proof. We use the notations from Section 6. Thus, $(R^\ell f)(1, 1) = (1, 1)_\ell$ and

$$(R^{\ell-1} f)(p, q) = (p, q)_{\ell-1} = A_{\ell-1}^{(p, q) \rightarrow (1, 1)} \cdot a$$

(by Theorem 7.6 (d), applied to $\ell - 1$ and (p, q) instead of ℓ and u). Solving this equation for $A_{\ell-1}^{(p, q) \rightarrow (1, 1)}$, we obtain

$$A_{\ell-1}^{(p, q) \rightarrow (1, 1)} = (R^{\ell-1} f)(p, q) \cdot \bar{a} \quad (27)$$

(since a is invertible). Note also that $R(R^{\ell-1} f) = R^\ell f \neq \perp$, and thus $(R^{\ell-1} f)(p, q)$ is invertible (by Lemma 3.22, applied to $R^{\ell-1} f$ and (p, q) instead of f and v).

Now,

$$\begin{aligned} (R^\ell f)(1, 1) &= (1, 1)_\ell = \overline{V_\ell^{(p, q) \rightarrow (1, 1)}} \cdot b && \text{(by Theorem 7.6 (c), applied to } u = (1, 1)\text{)} \\ &= \overline{A_{\ell-1}^{(p, q) \rightarrow (1, 1)}} \cdot b && \left(\text{since (23) yields } V_\ell^{(p, q) \rightarrow (1, 1)} = A_{\ell-1}^{(p, q) \rightarrow (1, 1)} \right) \\ &= \overline{\underbrace{(R^{\ell-1} f)(p, q) \cdot \bar{a}}_{= a \cdot \overline{(R^{\ell-1} f)(p, q)}}} \cdot b && \text{(by (27))} \\ & && \text{(since } (R^{\ell-1} f)(p, q) \text{ and } \bar{a} \text{ are invertible)} \\ &= a \cdot \overline{(R^{\ell-1} f)(p, q)} \cdot b. \end{aligned}$$

This proves Lemma 8.1. □

Unfortunately, our proof of Lemma 8.1 made use of the requirement that a be invertible, since $V_\ell^{(p, q) \rightarrow (1, 1)}$ and $A_{\ell-1}^{(p, q) \rightarrow (1, 1)}$ would not be well-defined otherwise. In order to remove this requirement, we make use of a trick, in which we “temporarily” set the label $f(0)$ to 1 and then argue that this has a predictable effect on $(Rf)(1, 1)$. This trick relies on the following:

Lemma 8.2. Let P be an arbitrary finite poset (not necessarily $[p] \times [q]$). Let $f, g \in \mathbb{K}^{\widehat{P}}$ be two \mathbb{K} -labellings such that $Rf \neq \perp$. Assume that

$$g(x) = f(x) \quad \text{for each } x \in \widehat{P} \setminus \{0\}. \quad (28)$$

Assume furthermore that $g(0) = 1$. Set $a = f(0)$. Then:

(a) We have $Rg \neq \perp$.

- (b) If $v \in P$ is not a minimal element of P , then $(Rf)(v) = (Rg)(v)$.
- (c) If $v \in P$ is a minimal element of P , then $(Rf)(v) = a \cdot (Rg)(v)$.

Proof of Lemma 8.2 (sketched). Our assumption (28) shows that the labels of f equal the corresponding labels of g at all elements of \widehat{P} other than at 0. Only the labels at 0 can differ.

Compute the labellings Rf and Rg recursively, as we did in Example 3.17, making sure to pick a linear extension of P that starts with all minimal elements of P (so that the toggles at these minimal elements all happen at the very end of our computation). The computation for Rf proceeds identically with the computation for Rg until we “interact with” the different labels at 0 – that is, until the labels $f(0)$ and $g(0)$ make an appearance in the sums $\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u)$ and $\sum_{\substack{u \in \widehat{P}; \\ u < v}} g(u)$, respectively (because all other labels of f equal the

corresponding labels of g). However, this “interaction” only happens when we toggle at a minimal element of P (since v has to be minimal in order for $f(0)$ to be an addend of the sum $\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u)$). Furthermore, when we do toggle at a minimal element v of P , the relevant sums $\sum_{\substack{u \in \widehat{P}; \\ u < v}} f(u)$ and $\sum_{\substack{u \in \widehat{P}; \\ u < v}} g(u)$ simplify to $f(0) = a$ and $g(0) = 1$, respectively

(because 0 is the **only** element $u \in \widehat{P}$ satisfying $u < v$). Therefore, the labels of Rf and Rg at v end up differing by a factor of a (more precisely, the value of Rf at v ends up being a times the label of Rg at v). This proves Lemma 8.2. \square

Let us now get rid of the “ a is invertible” requirement in Lemma 8.1:

Lemma 8.3. Assume that P is the $p \times q$ -rectangle $[p] \times [q]$. Let $\ell \in \mathbb{N}$ be such that $\ell \geq 1$. Let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^\ell f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Then,

$$(R^\ell f)(1, 1) = a \cdot \overline{(R^{\ell-1} f)(p, q)} \cdot b.$$

Proof. If $R^2 f \neq \perp$, then Lemma 3.24 yields that a and b are invertible (since $a = f(0)$ and $b = f(1)$), and therefore our claim follows directly from Lemma 8.1. For this reason, we WLOG assume that $R^2 f = \perp$. If we had $\ell \geq 2$, then we would thus conclude that $R^\ell f = \perp$ as well, which would contradict $R^\ell f \neq \perp$. Hence, we must have $\ell < 2$, so that $\ell = 1$. Therefore, $R^{\ell-1} = R^{1-1} = R^0 = \text{id}$ and consequently $(R^{\ell-1} f)(p, q) = f(p, q)$. Also, $R^\ell = R$ (since $\ell = 1$). Hence, $R = R^\ell$, so that $Rf = R^\ell f \neq \perp$.

Now, let $g \in \mathbb{K}^{\widehat{P}}$ be the \mathbb{K} -labelling that is obtained from f by replacing the label $f(0)$ by 1. Thus, we have

$$g(x) = f(x) \quad \text{for each } x \in \widehat{P} \setminus \{0\}, \quad (29)$$

and we have $g(0) = 1$. Then, Lemma 8.2 (a) yields $Rg \neq \perp$. In other words, $R^1 g \neq \perp$.

We have $(p, q) \in P \subseteq \widehat{P} \setminus \{0\}$. Hence, (29) yields $g(p, q) = f(p, q)$.

We also have $1 \in \widehat{P} \setminus \{0\}$. Thus, (29) yields $g(1) = f(1) = b$, so that $b = g(1)$. Also, $1 = g(0)$, and clearly 1 is invertible. Hence, Lemma 8.1 (applied to 1, g and 1 instead of ℓ , f and a) yields

$$(R^1 g)(1, 1) = 1 \cdot \overline{(R^{1-1} g)(p, q)} \cdot b = \overline{(R^{1-1} g)(p, q)} \cdot b.$$

In view of $R^1 = R$ and $R^{1-1} = \text{id}$, we can rewrite this as

$$(Rg)(1, 1) = \overline{g(p, q)} \cdot b.$$

However, $(1, 1)$ is a minimal element of P . Thus, Lemma 8.2 (c) (applied to $v = (1, 1)$) yields

$$(Rf)(1, 1) = a \cdot \underbrace{(Rg)(1, 1)}_{=\overline{g(p, q)} \cdot b} = a \cdot \overline{g(p, q)} \cdot b = a \cdot \overline{f(p, q)} \cdot b \quad (\text{since } g(p, q) = f(p, q)).$$

In view of $R^\ell = R$ and $(R^{\ell-1} f)(p, q) = f(p, q)$, we can rewrite this as

$$(R^\ell f)(1, 1) = a \cdot \overline{(R^{\ell-1} f)(p, q)} \cdot b.$$

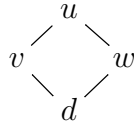
Thus, Lemma 8.3 is proven. \square

This settles the easiest case of Theorem 4.8 – namely, the case $(i, j) = (1, 1)$. To get a grip on the general case, we need more lemmas.

9. The conversion lemma

We continue using the notations from Section 6.

Lemma 9.1 (Four neighbors lemma). Let u, v, w, d be four adjacent elements of P that are arranged as follows on the Hasse diagram of P :



(i.e., we have $d = (i, j)$, $v = (i + 1, j)$, $w = (i, j + 1)$ and $u = (i + 1, j + 1)$ for some $i \in [p - 1]$ and some $j \in [q - 1]$).

Assume that a is invertible. Let $\ell \geq 1$ be such that $R^{\ell+1} f \neq \perp$. Then:

(a) We have

$$\overline{v}_\ell \cdot V_\ell^d \cdot d_\ell = \overline{u}_\ell \cdot A_\ell^u \cdot w_\ell.$$

(b) We have

$$\overline{w}_\ell \cdot V_\ell^d \cdot d_\ell = \overline{u}_\ell \cdot A_\ell^u \cdot v_\ell.$$

Proof. **(a)** We have $R(R^\ell f) = R^{\ell+1}f \neq \perp = R(\perp)$ and thus $R^\ell f \neq \perp$. Hence, Lemma 7.1 **(a)** yields that v_ℓ is invertible. Similarly, w_ℓ and u_ℓ and d_ℓ are invertible. Also, Lemma 7.1 **(d)** (applied to d instead of v) yields that the element V_ℓ^d is well-defined and invertible. Moreover, Lemma 7.1 **(c)** (applied to u and $\ell + 1$ instead of v and ℓ) yields that the element A_ℓ^u is well-defined and invertible.

The elements $s \in \widehat{P}$ that satisfy $s \succ d$ are v and w . Hence, $\sum_{s \succ d} \overline{s_\ell} = \overline{v_\ell} + \overline{w_\ell}$ (where, of course, the sum ranges over $s \in \widehat{P}$). Now, the definition of V_ℓ^d yields

$$V_\ell^d = \overline{\sum_{s \succ d} s_\ell} \cdot \overline{d_\ell} = \overline{v_\ell + w_\ell} \cdot \overline{d_\ell} \quad (30)$$

(since $\sum_{s \succ d} \overline{s_\ell} = \overline{v_\ell} + \overline{w_\ell}$).

The elements $s \in \widehat{P}$ that satisfy $s \prec u$ are v and w . Hence, $\sum_{s \prec u} s_\ell = v_\ell + w_\ell$. Now, the definition of A_ℓ^u yields

$$A_\ell^u = u_\ell \cdot \overline{\sum_{s \prec u} s_\ell} = u_\ell \cdot \overline{v_\ell + w_\ell} \quad (31)$$

(since $\sum_{s \prec u} s_\ell = v_\ell + w_\ell$). Since this is well-defined, the element $v_\ell + w_\ell$ of \mathbb{K} must be invertible. Also, we already know that v_ℓ and w_ℓ are invertible. Hence, Proposition 2.4 **(b)** (applied to v_ℓ and w_ℓ instead of a and b) yields that $\overline{v_\ell} + \overline{w_\ell}$ is invertible as well and its inverse is

$$\overline{v_\ell} + \overline{w_\ell} = v_\ell \cdot \overline{v_\ell + w_\ell} \cdot w_\ell.$$

Now,

$$\overline{v_\ell} \cdot \underbrace{V_\ell^d}_{=\overline{v_\ell + w_\ell} \cdot \overline{d_\ell} \text{ (by (30))}} \cdot d_\ell = \overline{v_\ell} \cdot \underbrace{\overline{v_\ell + w_\ell}}_{=v_\ell \cdot \overline{v_\ell + w_\ell} \cdot w_\ell} \cdot \underbrace{\overline{d_\ell} \cdot d_\ell}_{=1} = \underbrace{\overline{v_\ell} \cdot v_\ell}_{=1} \cdot \overline{v_\ell + w_\ell} \cdot w_\ell = \overline{v_\ell + w_\ell} \cdot w_\ell.$$

Comparing this with

$$\overline{u_\ell} \cdot \underbrace{A_\ell^u}_{=u_\ell \cdot \overline{v_\ell + w_\ell} \text{ (by (31))}} \cdot w_\ell = \underbrace{\overline{u_\ell} \cdot u_\ell}_{=1} \cdot \overline{v_\ell + w_\ell} \cdot w_\ell = \overline{v_\ell + w_\ell} \cdot w_\ell,$$

we obtain $\overline{v_\ell} \cdot V_\ell^d \cdot d_\ell = \overline{u_\ell} \cdot A_\ell^u \cdot w_\ell$. Thus, Lemma 9.1 **(a)** is proved.

(b) This can be proved by the same argument that we used to prove part **(a)** (with the roles of v and w interchanged). \square

We recall our conventions for drawing the $p \times q$ -rectangle $P = [p] \times [q]$. In light of these conventions, we shall refer to the set $\{(k, q) \mid k \in [p]\}$ as the *northeastern edge* of P , and to the set $\{(i, 1) \mid i \in [p]\}$ as the *southwestern edge* of P .

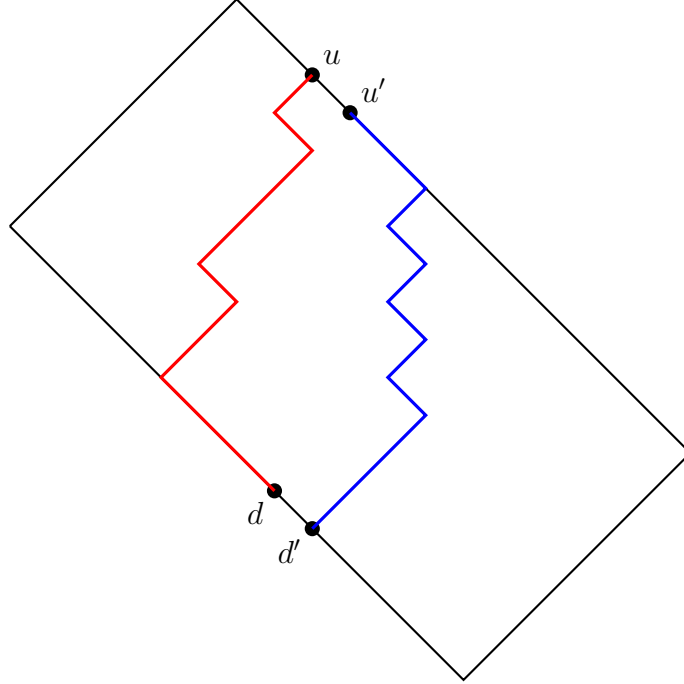
The next lemma is crucial, as it allows us to “convert” between A ’s and V ’s without changing the subscript.

Lemma 9.2 (Conversion lemma). Let u and u' be two elements of the northeastern edge of P satisfying $u \succ u'$ (that is, let $u = (k, q)$ and $u' = (k - 1, q)$ for some $k \in \{2, 3, \dots, p\}$). Let d and d' be two elements of the southwestern edge of P satisfying $d \succ d'$ (that is, let $d = (i, 1)$ and $d' = (i - 1, 1)$ for some $i \in \{2, 3, \dots, p\}$).

Assume that a is invertible. Let $\ell \geq 1$ be such that $R^{\ell+1}f \neq \perp$. Then we have:

$$A_\ell^{u \rightarrow d} = V_\ell^{u' \rightarrow d'}.$$

Here is an illustration for this lemma:



(the red path indexes one addend in the sum $A_\ell^{u \rightarrow d} = \sum_{\mathbf{p} \text{ is a path from } u \text{ to } d} A_\ell^{\mathbf{p}}$, while the blue path contributes to the sum $V_\ell^{u' \rightarrow d'} = \sum_{\mathbf{p} \text{ is a path from } u' \text{ to } d'} V_\ell^{\mathbf{p}}$).

In the case when \mathbb{K} is commutative, Lemma 9.2 was independently discovered by Johnson and Liu [JohLiu22]. More precisely, [JohLiu22, Lemma 4.1] extends it from sums over paths (such as $A_\ell^{u \rightarrow d}$ and $V_\ell^{u' \rightarrow d'}$) to sums over k -tuples of non-intersecting paths. It is unclear whether this extension can still be made when \mathbb{K} is not commutative (what order should the A_ℓ^v 's along different paths be multiplied in?), but the use of determinants likely precludes any noncommutative generalization of the proof in [JohLiu22].

Proof of Lemma 9.2. Let $\ell \in \mathbb{N}$. We “interpolate” between the paths from u to d and the paths from u' to d' using what we call “path-jump-paths”. To define these formally, we introduce some more basic notations.

The first coordinate of any $x \in P$ will be denoted by $\mathbf{first} x$. Thus, $\mathbf{first}(i, j) = i$ for any $(i, j) \in P$.

Furthermore, for any $x = (i, j) \in P$, we define the *rank* of x to be the positive integer $i + j - 1$. This rank will be denoted by $\mathbf{rank} x$.

We define a new binary relation \blacktriangleright on the set P as follows: If x and y are two elements of P , then the relation $x \blacktriangleright y$ means “ $\mathbf{rank} x = \mathbf{rank} y + 1$ and $\mathbf{first} x > \mathbf{first} y$ ”. In other words, the relation $x \blacktriangleright y$ means that

$$\text{if } x = (i, j), \text{ then } y = (i - k, j + k - 1) \text{ for some } k > 0.$$

Visually speaking, it means that y is one step southeast and a (nonnegative) amount of steps east of x (on the Hasse diagram).

We define a *path-jump-path* to be a tuple $\mathbf{p} = (v_0, v_1, \dots, v_k)$ of elements of P along with a chosen number $i \in \{0, 1, \dots, k - 1\}$ such that the chain of relations

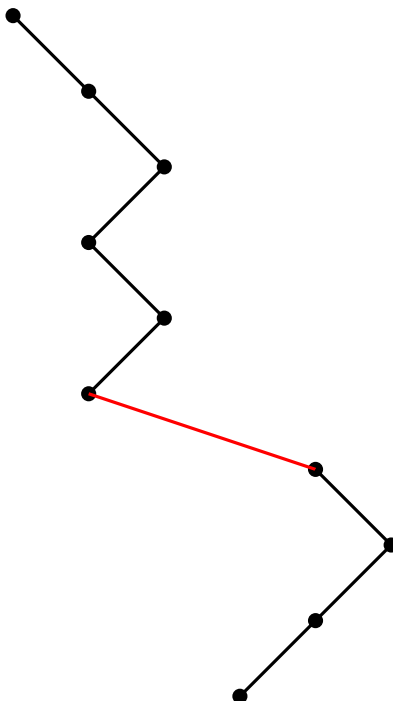
$$v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_k$$

holds. We denote this path-jump-path simply by

$$\mathbf{p} = (v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_k), \quad (32)$$

and we say that this path-jump-path \mathbf{p} has *jump at i* . The elements v_0, v_1, \dots, v_k are called the *vertices* of this path-jump-path. The pairs (v_j, v_{j+1}) of consecutive vertices are called the *steps* of this path-jump-path. Such a step (v_j, v_{j+1}) is said to be a \succ -step if $j \neq i$, and it is said to be a \blacktriangleright -step if $j = i$.

Here is an example of a path-jump-path, where the red edge is the \blacktriangleright -step:



(Note that two vertices x and y can satisfy $x \blacktriangleright y$ and $x \succ y$ simultaneously. Thus, it can happen that several path-jump-paths with jumps at different i 's contain the same

vertices. We nevertheless do not consider these path-jump-paths to be identical, because we understand a path-jump-path like (32) to “remember” not only its vertices v_0, v_1, \dots, v_k but also the value of i .)

A *path-jump-path from u to d'* will mean a path-jump-path $(v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_k)$ such that $v_0 = u$ and $v_k = d'$.

We note that if two elements x and y of P satisfy $x \succ y$ or $x \blacktriangleright y$, then

$$\mathbf{rank} y = \mathbf{rank} x - 1. \quad (33)$$

As a consequence of this fact, successive entries v_{j-1} and v_j in a path-jump-path $(v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_k)$ always satisfy $\mathbf{rank}(v_j) = \mathbf{rank}(v_{j-1}) - 1$ for each $j \in [k]$. In other words, the ranks of the vertices of a path-jump-path decrease by 1 at each step.

Hence, the difference in ranks between the first and final entries of a path-jump-path $(v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_k)$ is one less than its number of entries:

$$\mathbf{rank}(v_0) - \mathbf{rank}(v_k) = k. \quad (34)$$

Let $r := \mathbf{rank} u - \mathbf{rank}(d')$. Thus, any path-jump-path from u to d' must contain exactly $r + 1$ vertices (by (34)). In other words, any path-jump-path from u to d' must have the form $(v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_r)$.

We have $R(R^\ell f) = R^{\ell+1}f \neq \perp = R(\perp)$ and thus $R^\ell f \neq \perp$. Hence, Lemma 7.1 (a) yields that v_ℓ is well-defined and invertible for each $v \in P$. Also, Lemma 7.1 (d) yields that V_ℓ^v is well-defined and invertible for each $v \in P$. Moreover, Lemma 7.1 (c) (applied to $\ell + 1$ instead of ℓ) yields that A_ℓ^v is well-defined and invertible for each $v \in P$.

In this proof, we will not consider any \mathbb{K} -labellings other than $R^\ell f$. Thus, the only labels we will be using are the labels $v_\ell = (R^\ell f)(v)$ for $v \in \widehat{P}$. Thus, we agree to use the following shorthand notation: If $v \in \widehat{P}$, then the elements v_ℓ , V_ℓ^v and A_ℓ^v of \mathbb{K} will be denoted simply by v , V^v and A^v , respectively. In other words, **we shall omit subscripts when these subscripts are ℓ** . For instance, the product $A_\ell^u u_\ell \overline{u'_\ell}$ will thus be abbreviated as $A^u u \overline{u'}$.

For any path-jump-path

$$\mathbf{p} = (v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_r)$$

that contains $r + 1$ vertices, we set

$$E_{\mathbf{p}} := A^{v_0} A^{v_1} \dots A^{v_{i-1}} v_i \overline{v_{i+1}} V^{v_{i+2}} V^{v_{i+3}} \dots V^{v_r} \in \mathbb{K}.$$

Now we claim the following (again omitting subscripts that are ℓ):

Claim 1: We have

$$A^{u \rightarrow d} = \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } r-1}} E_{\mathbf{p}}.$$

Claim 2: We have

$$V^{u' \rightarrow d'} = \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } 0}} E_{\mathbf{p}}.$$

Claim 3: For each $j \in \{0, 1, \dots, r-2\}$, we have

$$\sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j}} E_{\mathbf{p}} = \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j+1}} E_{\mathbf{p}}.$$

Before we prove these three claims, let us explain how Lemma 9.2 will follow from them:

$$\begin{aligned} V_{\ell}^{u' \rightarrow d'} &= V^{u' \rightarrow d'} \\ &= \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } 0}} E_{\mathbf{p}} && (\text{by Claim 2}) \\ &= \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } 1}} E_{\mathbf{p}} && (\text{by Claim 3, applied to } j = 0) \\ &= \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } 2}} E_{\mathbf{p}} && (\text{by Claim 3, applied to } j = 1) \\ &= \dots \\ &= \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } r-1}} E_{\mathbf{p}} && (\text{by Claim 3, applied to } j = r-2) \\ &= A^{u \rightarrow d} && (\text{by Claim 1}) \\ &= A_{\ell}^{u \rightarrow d}. \end{aligned}$$

Hence, Lemma 9.2 will follow once Claims 1, 2 and 3 have been proved. Let us now prove these three claims:

Proof of Claim 1. We know that d lies on the southwestern edge of P . Hence, the only $s \in \widehat{P}$ satisfying $s \triangleleft d$ is d' (since $d \triangleright d'$). Therefore, $\sum_{\substack{s \in \widehat{P}; \\ s \triangleleft d}} s_{\ell} = d'_{\ell}$. However, the definition of A_{ℓ}^d shows that $A_{\ell}^d = d_{\ell} \cdot \overline{\sum_{\substack{s \in \widehat{P}; \\ s \triangleleft d}} s_{\ell}} = d_{\ell} \overline{d'_{\ell}}$ (since $\sum_{\substack{s \in \widehat{P}; \\ s \triangleleft d}} s_{\ell} = d'_{\ell}$). Since we omit subscripts (when these subscripts are ℓ), we can rewrite this as

$$A^d = d \overline{d'}. \tag{35}$$

We know that any path-jump-path from u to d' must have the form $(v_0 \succ v_1 \succ \cdots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \cdots \succ v_r)$. If such a path-jump-path has jump at $r-1$, then it must have the form $(v_0 \succ v_1 \succ \cdots \succ v_{r-1} \blacktriangleright v_r)$; that is, its last step (v_{r-1}, v_r) is an \blacktriangleright -step. However, since it ends at d' , we must have $v_r = d'$ and thus $v_{r-1} \blacktriangleright v_r = d'$. This entails $v_{r-1} = d$ (since the only $g \in P$ satisfying $g \blacktriangleright d'$ is d ¹⁸), and therefore $(v_{r-1}, v_r) = (d, d')$ (since $v_r = d'$). In other words, the last step of this path-jump-path is (d, d') .

We have thus shown that if a path-jump-path from u to d' has jump at $r-1$, then its last step is (d, d') . Hence, any path-jump-path from u to d' with jump at $r-1$ must have the form

$$(v_0 \succ v_1 \succ \cdots \succ v_{r-1} \blacktriangleright d'),$$

where $(v_0 \succ v_1 \succ \cdots \succ v_{r-1})$ is a path from u to d . Conversely, any tuple of the latter form is a path-jump-path from u to d' with jump at $r-1$ (since $d \blacktriangleright d'$). Therefore,

$$\begin{aligned} \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } r-1}} E_{\mathbf{p}} &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{r-1}) \\ \text{is a path from } u \text{ to } d}} \underbrace{E_{(v_0 \succ v_1 \succ \cdots \succ v_{r-1} \blacktriangleright d')}}_{=A^{v_0} A^{v_1} \cdots A^{v_{r-2}} v_{r-1} \bar{d}'} \\ &\quad \text{(by the definition of } E_{(v_0 \succ v_1 \succ \cdots \succ v_{r-1} \blacktriangleright d')}) \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{r-1}) \\ \text{is a path from } u \text{ to } d}} A^{v_0} A^{v_1} \cdots A^{v_{r-2}} \underbrace{v_{r-1} \bar{d}'}_{=d} \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{r-1}) \\ \text{is a path from } u \text{ to } d}} A^{v_0} A^{v_1} \cdots A^{v_{r-2}} \underbrace{\bar{d}d'}_{=A^d \text{ (by (35))}} \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{r-1}) \\ \text{is a path from } u \text{ to } d}} A^{v_0} A^{v_1} \cdots A^{v_{r-2}} \underbrace{A^d}_{=A^{v_{r-1}} \text{ (because } d=v_{r-1})} \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{r-1}) \\ \text{is a path from } u \text{ to } d}} \underbrace{A^{v_0} A^{v_1} \cdots A^{v_{r-2}} A^{v_{r-1}}}_{=A^{(v_0 \succ v_1 \succ \cdots \succ v_{r-1})}} \\ &\quad \text{(by the definition of } A^{(v_0 \succ v_1 \succ \cdots \succ v_{r-1})}) \\ &= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{r-1}) \\ \text{is a path from } u \text{ to } d}} A^{(v_0 \succ v_1 \succ \cdots \succ v_{r-1})} = \sum_{\mathbf{p} \text{ is a path from } u \text{ to } d} A^{\mathbf{p}} \\ &= A^{u \rightarrow d} \quad \text{(by the definition of } A^{u \rightarrow d}). \end{aligned}$$

This proves Claim 1. □

Proof of Claim 2. This is analogous to the proof of Claim 1. This time, we need to argue that if a path-jump-path from u to d' has jump at 0, then its first step is (u, u') (since the only $g \in P$ satisfying $u \blacktriangleright g$ is u'). □

Proving Claim 3 is a bit trickier. As an auxiliary result, we first show the following:

¹⁸This follows easily from the geographical positions of d and d' on the southwestern edge of P .

Claim 4: Let s and t be two elements of P . Then,

$$\sum_{\substack{x \in P; \\ s \blacktriangleright x \triangleright t}} s \bar{x} V^t = \sum_{\substack{x \in P; \\ s \triangleright x \blacktriangleright t}} A^s x \bar{t}. \quad (36)$$

Proof of Claim 4. First, we observe that an $x \in P$ satisfying $s \blacktriangleright x \triangleright t$ cannot exist unless $\mathbf{rank} t = \mathbf{rank} s - 2$ (because (33) yields that such an x must satisfy $\mathbf{rank} x = \mathbf{rank} s - 1$ and $\mathbf{rank} t = \mathbf{rank} x - 1$, whence $\mathbf{rank} t = \mathbf{rank} x - 1 = (\mathbf{rank} s - 1) - 1 = \mathbf{rank} s - 2$). Hence, the left hand side of the desired equality (36) is an empty sum unless $\mathbf{rank} t = \mathbf{rank} s - 2$. Similarly, the same can be said about the right hand side. Thus, (36) boils down to $0 = 0$ unless $\mathbf{rank} t = \mathbf{rank} s - 2$. We therefore assume WLOG that $\mathbf{rank} t = \mathbf{rank} s - 2$. In other words, $\mathbf{rank} s - \mathbf{rank} t = 2$.

The definition of V_ℓ^t yields $V_\ell^t = \sum_{x \triangleright t} \overline{x_\ell} \cdot \bar{t}_\ell$. Omitting the subscripts, we can rewrite this as

$$V^t = \sum_{x \triangleright t} \overline{x} \cdot \bar{t}. \quad (37)$$

The definition of A_ℓ^s yields $A_\ell^s = s_\ell \cdot \sum_{x \triangleleft s} \overline{x_\ell}$. Omitting the subscripts, we can rewrite this as

$$A^s = s \cdot \sum_{x \triangleleft s} \overline{x}. \quad (38)$$

Write $s \in P$ in the form $s = (i, j)$. Write $t \in P$ in the form $t = (i', j')$. From $s = (i, j)$, we obtain $\mathbf{rank} s = i + j - 1$. Likewise, $\mathbf{rank} t = i' + j' - 1$. Hence,

$$\mathbf{rank} s - \mathbf{rank} t = (i + j - 1) - (i' + j' - 1) = i + j - i' - j',$$

so that $i + j - i' - j' = \mathbf{rank} s - \mathbf{rank} t = 2$. Thus, $j' = i + j - i' - 2$.

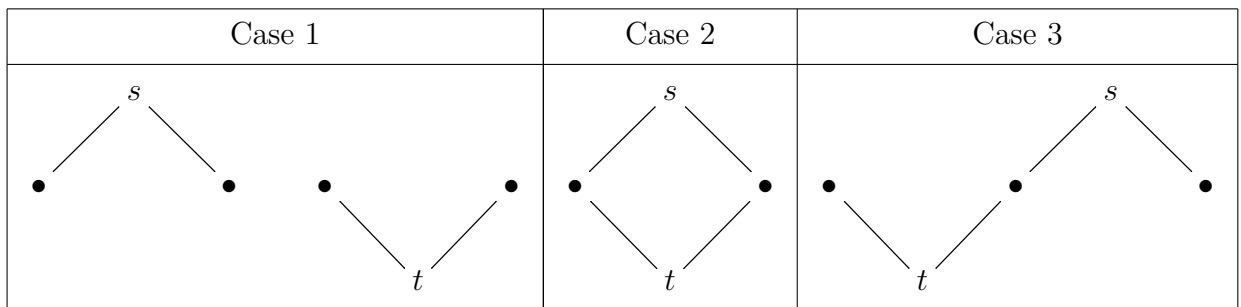
We are in one of the following three cases:

Case 1: We have $i' < i - 1$.

Case 2: We have $i' = i - 1$.

Case 3: We have $i' > i - 1$.

Representative examples for these three cases are illustrated in the following pictures:



(the bullets signify the positions of potential neighbors of s and t ; some of these positions may fall outside of P , but this does not disturb our argument).

Let us first consider Case 1. In this case, we have $i' < i - 1$. Thus, $i' + 1 < i$. Hence, each element x of P that satisfies $x \succ t$ must satisfy $s \blacktriangleright x$ automatically¹⁹. Therefore, the summation sign $\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}}$ can be simplified to $\sum_{\substack{x \in P; \\ x \succ t}}$. In turn, the latter summation sign $\sum_{\substack{x \in P; \\ x \succ t}}$ can be rewritten as $\sum_{\substack{x \in \widehat{P}; \\ x \succ t}}$ (because any $x \in \widehat{P}$ that satisfies $x \succ t$ must automatically belong to P ²⁰). Hence, we obtain the following equality of summation signs:

$$\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} = \sum_{\substack{x \in P; \\ x \succ t}} = \sum_{\substack{x \in \widehat{P}; \\ x \succ t}} = \sum_{x \succ t}$$

(since our sums are understood to range over \widehat{P} by default). Thus,

$$\begin{aligned} \sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s \bar{x} V^t &= \sum_{x \succ t} s \bar{x} V^t = s \left(\sum_{x \succ t} \bar{x} \right) \underbrace{V^t}_{\substack{= \sum_{x \succ t} \bar{x} \cdot \bar{t} \\ \text{(by (37))}}} = s \underbrace{\left(\sum_{x \succ t} \bar{x} \right) \sum_{x \succ t} \bar{x} \cdot \bar{t}}_{=1} \\ &= s \bar{t}. \end{aligned} \tag{39}$$

Furthermore, $i' < i - 1$, so that $i - 1 > i'$. Hence, each element x of P that satisfies $s \succ x$ must satisfy $x \blacktriangleright t$ automatically²¹. Therefore, the summation sign $\sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}}$

¹⁹*Proof.* Let $x \in P$ be such that $x \succ t$. We must show that $s \blacktriangleright x$.

From $x \succ t$, we obtain $\text{rank } t = \text{rank } x - 1$ (by (33), applied to $y = t$) and thus $\text{rank } x - \text{rank } t = 1$. Now,

$$\text{rank } s - \text{rank } x = \underbrace{(\text{rank } s - \text{rank } t)}_{=2} - \underbrace{(\text{rank } x - \text{rank } t)}_{=1} = 2 - 1 = 1.$$

Thus, $\text{rank } s = \text{rank } x + 1$.

Furthermore, x is either $(i' + 1, j')$ or $(i', j' + 1)$ (since $x \succ t = (i', j')$). Hence, $\text{first } x \leq i' + 1 < i = \text{first } s$ (since $s = (i, j)$), so that $\text{first } s > \text{first } x$. Combining this with $\text{rank } s = \text{rank } x + 1$, we obtain $s \blacktriangleright x$ (by the definition of the relation \blacktriangleright).

²⁰*Proof.* From $t = (i', j')$, we obtain $\text{first } t = i' < i - 1 < i \leq p = \text{first } (p, q)$. Hence, $\text{first } t \neq \text{first } (p, q)$, so that $t \neq (p, q)$. This shows that t is not a maximal element of P (since the only maximal element of P is (p, q)). In other words, we don't have $t \leq 1$ in \widehat{P} . In other words, we don't have $1 \succ t$ in \widehat{P} . Hence, any $x \in \widehat{P}$ that satisfies $x \succ t$ must automatically belong to P (since we have neither $1 \succ t$ nor $0 \succ t$).

²¹*Proof.* Let $x \in P$ be such that $s \succ x$. We must show that $x \blacktriangleright t$.

From $s \succ x$, we obtain $\text{rank } x = \text{rank } s - 1$ (by (33), applied to s and x instead of x and y). Hence, $\text{rank } s - \text{rank } x = 1$, so that

$$\text{rank } x - \text{rank } t = \underbrace{(\text{rank } s - \text{rank } t)}_{=2} - \underbrace{(\text{rank } s - \text{rank } x)}_{=1} = 2 - 1 = 1.$$

Thus, $\text{rank } x = \text{rank } t + 1$.

Furthermore, x is either $(i - 1, j)$ or $(i, j - 1)$ (since $(i, j) = s \succ x$). Hence, $\text{first } x \geq i - 1 > i' = \text{first } t$ (since $t = (i', j')$). Combining this with $\text{rank } x = \text{rank } t + 1$, we obtain $x \blacktriangleright t$ (by the definition of the relation \blacktriangleright).

to $\sum_{\substack{x \in P; \\ s \succ x}}$. In turn, the latter summation sign $\sum_{\substack{x \in P; \\ s \succ x}}$ can be rewritten as $\sum_{\substack{x \in \widehat{P}; \\ s \succ x}}$ (because any $x \in \widehat{P}$ that satisfies $s \succ x$ must automatically belong to P ²²). Hence, we obtain the following equality of summation signs:

$$\sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} = \sum_{\substack{x \in P; \\ s \succ x}} = \sum_{\substack{x \in \widehat{P}; \\ s \succ x}} = \sum_{\substack{x \in \widehat{P}; \\ x < s}} = \sum_{x < s}$$

(since our sums are understood to range over \widehat{P} by default). Thus,

$$\sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} A^s x \bar{t} = \sum_{x < s} A^s x \bar{t} = \underbrace{A^s}_{=s \cdot \sum_{\substack{x \in P; \\ x < s}} x}_{\text{(by (38))}} \left(\sum_{x < s} x \right) \bar{t} = s \cdot \underbrace{\sum_{x < s} x}_{=1} \left(\sum_{x < s} x \right) \bar{t} = s \bar{t}.$$

Comparing this with (39), we obtain $\sum_{\substack{x \in P; \\ s \blacktriangleright x \blacktriangleright t}} s \bar{x} V^t = \sum_{\substack{x \in P; \\ s \blacktriangleright x \blacktriangleright t}} A^s x \bar{t}$. Thus, Claim 4 is proved in

Case 1.

Let us now consider Case 2. In this case, we have $i' = i - 1$. Hence, $j' = i + j - \underbrace{i'}_{=i-1} - 2 = i + j - (i - 1) - 2 = j - 1$. Thus, $t = (i', j') = (i - 1, j - 1)$ (since $i' = i - 1$ and $j' = j - 1$). Let $v := (i, j - 1)$ and $w := (i - 1, j)$. In our coordinate system, the four points

$$s = (i, j), \quad t = (i - 1, j - 1), \quad v = (i, j - 1), \quad w = (i - 1, j)$$

are arranged in a 1×1 -square, which looks as follows:

$$\begin{array}{ccc} & s & \\ & \diagdown \quad \diagup & \\ v & & w \\ & \diagup \quad \diagdown & \\ & t & \end{array} \quad . \quad (40)$$

Hence, v and w belong to P (since s and t belong to P), and furthermore, Lemma 9.1 (b) (applied to s and t instead of u and d) yields

$$\bar{w}_\ell \cdot V_\ell^t \cdot t_\ell = \bar{s}_\ell \cdot A_\ell^s \cdot v_\ell.$$

Since we are omitting subscripts, we can rewrite this as follows:

$$\bar{w} \cdot V^t \cdot t = \bar{s} \cdot A^s \cdot v.$$

²²*Proof.* From $s = (i, j)$, we obtain $\mathbf{first} s = i > i - 1 > i' \geq 1 = \mathbf{first}(1, 1)$. Hence, $\mathbf{first} s \neq \mathbf{first}(1, 1)$, so that $s \neq (1, 1)$. This shows that s is not a minimal element of P (since the only minimal element of P is $(1, 1)$). In other words, we don't have $s \succ 0$ in \widehat{P} . Hence, any $x \in \widehat{P}$ that satisfies $s \succ x$ must automatically belong to P (since we have neither $s \succ 0$ nor $s \succ 1$).

The picture (40) shows that we have $s \blacktriangleright w$ but not $s \blacktriangleright v$. Hence, there is only one element $x \in P$ that satisfies $s \blacktriangleright x \succ t$; namely, this element x is w . Hence,

$$\begin{aligned} \sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s\bar{x}V^t &= s\bar{w}V^t = s \cdot \bar{w} \cdot V^t \cdot \underbrace{1}_{=t\bar{t}} = s \cdot \underbrace{\bar{w} \cdot V^t \cdot t}_{=\bar{s} \cdot A^s \cdot v} \cdot \bar{t} = \underbrace{s \cdot \bar{s}}_{=1} \cdot A^s \cdot v \cdot \bar{t} \\ &= A^s \cdot v \cdot \bar{t}. \end{aligned} \tag{41}$$

On the other hand, the picture (40) shows that we have $v \blacktriangleright t$ but not $w \blacktriangleright t$. Hence, there is only one element $x \in P$ that satisfies $s \succ x \blacktriangleright t$; namely, this element x is v . Hence,

$$\sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} A^s x\bar{t} = A^s v\bar{t} = A^s \cdot v \cdot \bar{t}.$$

Comparing this with (41), we obtain $\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s\bar{x}V^t = \sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} A^s x\bar{t}$. Thus, Claim 4 is proved in

Case 2.

Let us finally consider Case 3. In this case, we have $i' > i - 1$. Thus, $i' \geq i$ (since i' and i are integers), so that $i \leq i'$. Note that $i = \mathbf{first} s$ (since $s = (i, j)$) and $i' = \mathbf{first} t$ (since $t = (i', j')$).

There exists no $x \in P$ satisfying $s \blacktriangleright x \succ t$ (because if $x \in P$ satisfies $s \blacktriangleright x \succ t$, then $x \succ t = (i', j')$ entails $\mathbf{first} x \geq i' \geq i = \mathbf{first} s$, but this clearly contradicts $s \blacktriangleright x$). Hence, the sum $\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s\bar{x}V^t$ is empty. Thus, $\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s\bar{x}V^t = 0$.

Furthermore, there exists no $x \in P$ satisfying $s \succ x \blacktriangleright t$ (because if $x \in P$ satisfies $s \succ x \blacktriangleright t$, then $(i, j) = s \succ x$ entails $\mathbf{first} x \leq i \leq i' = \mathbf{first} t$; but this clearly contradicts $x \blacktriangleright t$). Hence, the sum $\sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} A^s x\bar{t}$ is empty. Thus, $\sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} A^s x\bar{t} = 0$.

Comparing this with $\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s\bar{x}V^t = 0$, we obtain $\sum_{\substack{x \in P; \\ s \blacktriangleright x \succ t}} s\bar{x}V^t = \sum_{\substack{x \in P; \\ s \succ x \blacktriangleright t}} A^s x\bar{t}$. Thus, Claim 4 is proved in Case 3.

We have now proved Claim 4 in all three cases. \square

We can now step to the proof of Claim 3:

Proof of Claim 3. Let $j \in \{0, 1, \dots, r - 2\}$.

We know that any path-jump-path from u to d' must have the form $(v_0 \succ v_1 \succ \dots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \dots \succ v_r)$. If such a path-jump-path has jump at j , then

it must have the form $(v_0 \succ v_1 \succ \cdots \succ v_j \blacktriangleright v_{j+1} \succ v_{j+2} \succ \cdots \succ v_r)$. Thus,

$$\begin{aligned}
& \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j}} E_{\mathbf{p}} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j \blacktriangleright v_{j+1} \succ v_{j+2} \succ \cdots \succ v_r) \\ \text{is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j}} \underbrace{E_{(v_0 \succ v_1 \succ \cdots \succ v_j \blacktriangleright v_{j+1} \succ v_{j+2} \succ \cdots \succ v_r)}}_{= A^{v_0} A^{v_1} \cdots A^{v_{j-1}} v_j \overline{v_{j+1}} V^{v_{j+2}} V^{v_{j+3}} \cdots V^{v_r}} \\
&\quad \text{(by the definition of } E_{(v_0 \succ v_1 \succ \cdots \succ v_j \blacktriangleright v_{j+1} \succ v_{j+2} \succ \cdots \succ v_r)}) \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j \blacktriangleright v_{j+1} \succ v_{j+2} \succ \cdots \succ v_r) \\ \text{is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j}} A^{v_0} A^{v_1} \cdots A^{v_{j-1}} v_j \overline{v_{j+1}} V^{v_{j+2}} V^{v_{j+3}} \cdots V^{v_r} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} \sum_{\substack{v_{j+1} \in P; \\ v_j \blacktriangleright v_{j+1} \succ v_{j+2}}} A^{v_0} A^{v_1} \cdots A^{v_{j-1}} v_j \overline{v_{j+1}} V^{v_{j+2}} V^{v_{j+3}} \cdots V^{v_r} \\
&\quad \left(\begin{array}{c} \text{here, we have broken up our} \\ \text{path-jump-path } (v_0 \succ v_1 \succ \cdots \succ v_j \blacktriangleright v_{j+1} \succ v_{j+2} \succ \cdots \succ v_r) \\ \text{into two paths } (v_0 \succ v_1 \succ \cdots \succ v_j) \text{ and } (v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{and an intermediate vertex } v_{j+1} \text{ satisfying } v_j \blacktriangleright v_{j+1} \succ v_{j+2} \end{array} \right) \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} \sum_{\substack{x \in P; \\ v_j \blacktriangleright x \succ v_{j+2}}} A^{v_0} A^{v_1} \cdots A^{v_{j-1}} v_j \overline{x} V^{v_{j+2}} V^{v_{j+3}} \cdots V^{v_r} \\
&\quad \text{(here we have renamed } v_{j+1} \text{ as } x \text{ in the inner sum)} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} A^{v_0} A^{v_1} \cdots A^{v_{j-1}} \underbrace{\sum_{\substack{x \in P; \\ v_j \blacktriangleright x \succ v_{j+2}}} v_j \overline{x} V^{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r}}_{= \sum_{\substack{x \in P; \\ v_j \blacktriangleright x \succ v_{j+2}}} A^{v_j} x \overline{v_{j+2}}} \\
&\quad \text{(by Claim 4, applied to } s=v_j \text{ and } t=v_{j+2}) \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} A^{v_0} A^{v_1} \cdots A^{v_{j-1}} \sum_{\substack{x \in P; \\ v_j \blacktriangleright x \succ v_{j+2}}} A^{v_j} x \overline{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r}.
\end{aligned}$$

We know that any path-jump-path from u to d' must have the form $(v_0 \succ v_1 \succ \cdots \succ v_i \blacktriangleright v_{i+1} \succ v_{i+2} \succ \cdots \succ v_r)$. If such a path-jump-path has jump at $j+1$,

then it must have the form $(v_0 \succ v_1 \succ \cdots \succ v_{j+1} \blacktriangleright v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r)$. Thus,

$$\begin{aligned}
& \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j+1}} E_{\mathbf{p}} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{j+1} \blacktriangleright v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j+1}} \underbrace{E_{(v_0 \succ v_1 \succ \cdots \succ v_{j+1} \blacktriangleright v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r)}}_{= A^{v_0} A^{v_1} \cdots A^{v_j} v_{j+1} \overline{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r}} \\
&\quad \text{(by the definition of } E_{(v_0 \succ v_1 \succ \cdots \succ v_{j+1} \blacktriangleright v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r)}) \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_{j+1} \blacktriangleright v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j+1}} A^{v_0} A^{v_1} \cdots A^{v_j} v_{j+1} \overline{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} \sum_{\substack{v_{j+1} \in P; \\ v_j \succ v_{j+1} \blacktriangleright v_{j+2}}} A^{v_0} A^{v_1} \cdots A^{v_j} v_{j+1} \overline{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} \sum_{\substack{x \in P; \\ v_j \succ x \blacktriangleright v_{j+2}}} A^{v_0} A^{v_1} \cdots A^{v_j} x \overline{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r} \\
&\quad \text{(here we have renamed } v_{j+1} \text{ as } x \text{ in the inner sum)} \\
&= \sum_{\substack{(v_0 \succ v_1 \succ \cdots \succ v_j) \\ \text{is a path starting at } u}} \sum_{\substack{(v_{j+2} \succ v_{j+3} \succ \cdots \succ v_r) \\ \text{is a path ending at } d'}} A^{v_0} A^{v_1} \cdots A^{v_{j-1}} \sum_{\substack{x \in P; \\ v_j \succ x \blacktriangleright v_{j+2}}} A^{v_j} x \overline{v_{j+2}} V^{v_{j+3}} V^{v_{j+4}} \cdots V^{v_r}.
\end{aligned}$$

Comparing our last two equalities, we obtain

$$\sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j}} E_{\mathbf{p}} = \sum_{\substack{\mathbf{p} \text{ is a path-jump-path} \\ \text{from } u \text{ to } d' \\ \text{with jump at } j+1}} E_{\mathbf{p}}.$$

Thus, Claim 3 is proven. \square

We have now proved all three Claims 1, 2 and 3. As we explained, this completes the proof of Lemma 9.2. \square

Remark 9.3. Parts of the above proof of Lemma 9.2 can be rewritten in a more abstract manner, avoiding the notion of a “path-jump-path” and the nested sums that appeared in our proof of Claim 3.

To rewrite the proof, we need the notion of $P \times P$ -matrices. A $P \times P$ -matrix is a matrix whose rows and columns are indexed not by integers but by elements of P . (That is, it is a family of elements of \mathbb{K} indexed by pairs $(i, j) \in P \times P$.) If C is any $P \times P$ -matrix, and if i and j are two elements of P , then the (i, j) -th entry of C is denoted by $C_{i,j}$. Addition and multiplication are defined for $P \times P$ -matrices in the

same way as they are for usual matrices. That is, for any $P \times P$ -matrices C and D and any $(i, j) \in P \times P$, we have

$$(C + D)_{i,j} = C_{i,j} + D_{i,j} \quad \text{and} \quad (CD)_{i,j} = \sum_{k \in P} C_{i,k} D_{k,j}.$$

For any statement \mathcal{A} , we let $[\mathcal{A}]$ be the Iverson bracket (i.e., truth value) of \mathcal{A} . That is, $[\mathcal{A}] = 1$ if \mathcal{A} is true, and $[\mathcal{A}] = 0$ if \mathcal{A} is false.

Now, let $\ell \in \mathbb{N}$. Define three $P \times P$ -matrices \mathbf{A} , \mathbf{V} and \mathbf{U} by

$$\begin{aligned} \mathbf{A}_{x,y} &:= A^x [x \succ y], \\ \mathbf{V}_{x,y} &:= V^y [x \succ y], \\ \mathbf{U}_{x,y} &:= x\bar{y} [x \blacktriangleright y] \quad \text{for all } x, y \in P. \end{aligned}$$

Here, the relation $x \blacktriangleright y$ is defined as in the above proof of Lemma 9.2, and we are again omitting the “ ℓ ” subscripts, so (for instance) “ $x\bar{y}$ ” actually means $x_{\ell}\bar{y}_{\ell}$.

Now, Claim 4 in our above proof of Lemma 9.2 can be rewritten in a nice and compact form as the equality

$$\mathbf{A}\mathbf{U} = \mathbf{U}\mathbf{V}.$$

From this, we easily obtain

$$\mathbf{A}^k \mathbf{U} = \mathbf{U} \mathbf{V}^k \quad \text{for any } k \in \mathbb{N}. \quad (42)$$

This equality essentially replaces Claim 3 in the above proof.

Setting $k = \mathbf{ran} \mathfrak{k} u - \mathbf{ran} \mathfrak{k} d$ in (42), and comparing the (u, d') -entries of both sides, we quickly obtain $A^{u \rightarrow d} = V^{u' \rightarrow d'}$ (since $x \blacktriangleright d'$ holds only for $x = d$, and since $u \blacktriangleright x$ holds only for $x = u'$). This proves Lemma 9.2 again.

10. Proof of reciprocity: the case $j = 1$

Using the conversion lemma, we can now easily prove Theorem 4.8 in the case when $j = 1$:

Lemma 10.1. Assume that P is the $p \times q$ -rectangle $[p] \times [q]$. Let $i \in [p]$. Let $\ell \in \mathbb{N}$ satisfy $\ell \geq i$. Let $f \in \mathbb{K}^{\hat{P}}$ be a \mathbb{K} -labelling such that $R^{\ell} f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Then, using the notations from Section 6, we have

$$(i, 1)_{\ell} = a \cdot \overline{(p+1-i, q)_{\ell-i}} \cdot b.$$

Proof. We have $\ell \geq i \geq 1$. Hence, Lemma 8.3 yields that $(R^{\ell} f)(1, 1) = a \cdot \overline{(R^{\ell-1} f)(p, q)} \cdot b$. In other words, $(1, 1)_{\ell} = a \cdot \overline{(p, q)_{\ell-1}} \cdot b$. This proves Lemma 10.1 in the case when $i = 1$.

Hence, for the rest of this proof, we WLOG assume that $i \neq 1$. Thus, $i \geq 2$, so that $\ell \geq i \geq 2$, and therefore $R^2 f \neq \perp$ (by Lemma 3.21, since $R^{\ell} f \neq \perp$). Hence, Lemma 3.24

yields that a and b are invertible (since $a = f(0)$ and $b = f(1)$).

We have $\ell - i + 1 \geq 1$ (since $\ell \geq i$) and $\ell - i + 1 \leq \ell$ (since $i \geq 1$). The latter inequality entails $R^{\ell-i+1}f \neq \perp$ (by Lemma 3.21, since $R^\ell f \neq \perp$). Thus, Lemma 7.1 **(b)** (applied to $(p - i + 1, q)$ and $\ell - i + 1$ instead of v and ℓ) yields that the element $(p - i + 1, q)_{\ell-i}$ is well-defined and invertible.

Theorem 7.6 **(d)** (applied to $(p - i + 1, q)$ and $\ell - i$ instead of u and ℓ) yields

$$(p - i + 1, q)_{\ell-i} = A_{\ell-i}^{(p-i+1, q) \rightarrow (1, 1)} \cdot a.$$

Solving this for $A_{\ell-i}^{(p-i+1, q) \rightarrow (1, 1)}$, we obtain $A_{\ell-i}^{(p-i+1, q) \rightarrow (1, 1)} = (p - i + 1, q)_{\ell-i} \cdot \bar{a}$, and thus

$$\overline{A_{\ell-i}^{(p-i+1, q) \rightarrow (1, 1)}} = \overline{(p - i + 1, q)_{\ell-i} \cdot \bar{a}} = a \cdot \overline{(p - i + 1, q)_{\ell-i}}. \quad (43)$$

For each $k \in \{0, 1, \dots, i - 2\}$, we have

$$\begin{aligned} V_{\ell-k}^{(p-k, q) \rightarrow (i-k, 1)} &= A_{\ell-k-1}^{(p-k, q) \rightarrow (i-k, 1)} && \left(\begin{array}{l} \text{by (23), since we can easily} \\ \text{find } \ell - k \geq 2 \geq 1 \text{ and } R^{\ell-k}f \neq \perp \end{array} \right) \\ &= V_{\ell-k-1}^{(p-k-1, q) \rightarrow (i-k-1, 1)} && \left(\begin{array}{l} \text{by Lemma 9.2, applied to} \\ (p - k, q), (p - k - 1, q), \\ (i - k, 1), (i - k - 1, 1) \\ \text{and } \ell - k - 1 \\ \text{instead of } u, u', d, d' \text{ and } \ell \end{array} \right) \\ &= V_{\ell-(k+1)}^{(p-(k+1), q) \rightarrow (i-(k+1), 1)}. && (44) \end{aligned}$$

Now,

$$\begin{aligned} V_{\ell}^{(p, q) \rightarrow (i, 1)} &= V_{\ell-0}^{(p-0, q) \rightarrow (i-0, 1)} \\ &= V_{\ell-1}^{(p-1, q) \rightarrow (i-1, 1)} && \text{(by (44), applied to } k = 0) \\ &= V_{\ell-2}^{(p-2, q) \rightarrow (i-2, 1)} && \text{(by (44), applied to } k = 1) \\ &= \dots \\ &= V_{\ell-(i-1)}^{(p-(i-1), q) \rightarrow (i-(i-1), 1)} && \text{(by (44), applied to } k = i - 2) \\ &= V_{\ell-i+1}^{(p-i+1, q) \rightarrow (1, 1)} && \left(\begin{array}{l} \text{since } p - (i - 1) = p - i + 1 \\ \text{and } i - (i - 1) = 1 \\ \text{and } \ell - (i - 1) = \ell - i + 1 \end{array} \right) \\ &= A_{\ell-i}^{(p-i+1, q) \rightarrow (1, 1)} && \text{(by (23)).} \end{aligned} \quad (45)$$

However, Theorem 7.6 **(c)** (applied to $u = (i, 1)$) yields

$$\begin{aligned} (i, 1)_{\ell} &= \overline{V_{\ell}^{(p, q) \rightarrow (i, 1)}} \cdot b = \overline{A_{\ell-i}^{(p-i+1, q) \rightarrow (1, 1)}} \cdot b && \text{(by (45))} \\ &= a \cdot \overline{(p - i + 1, q)_{\ell-i}} \cdot b && \text{(by (43))} \\ &= a \cdot \overline{(p + 1 - i, q)_{\ell-i}} \cdot b. \end{aligned}$$

This proves Lemma 10.1. □

In analogy to Lemma 10.1, we have the following:

Lemma 10.2. Assume that P is the $p \times q$ -rectangle $[p] \times [q]$. Let $j \in [q]$. Let $\ell \in \mathbb{N}$ satisfy $\ell \geq j$. Let $f \in \mathbb{K}^{\hat{P}}$ be a \mathbb{K} -labelling such that $R^\ell f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Then, using the notations from Section 6, we have

$$(1, j)_\ell = a \cdot \overline{(p, q+1-j)_{\ell-j}} \cdot b.$$

Proof. The two coordinates u and v of an element $(u, v) \in P$ play symmetric roles. Lemma 10.2 is just Lemma 10.1 with the roles of these two coordinates interchanged. Thus, the proof of Lemma 10.2 is analogous to the proof of Lemma 10.1. \square

11. Proof of reciprocity: the general case

Somewhat surprisingly, the general case of Theorem 4.8 follows by a fairly straightforward induction argument from Lemma 10.1:

Proof of Theorem 4.8. We again use the notations from Section 6.

For any $(i, j) \in P$, we define $\mathbf{tilt}(i, j)$ to be the positive integer $i + 2j$. Our goal is to prove (10) for each $x = (i, j) \in P$ and $\ell \in \mathbb{N}$ satisfying $\ell - i - j + 1 \geq 0$ and $R^\ell f \neq \perp$. We will now prove this by strong induction on $\mathbf{tilt} x$.

Induction step: Fix $N \in \mathbb{N}$. Assume (as the induction hypothesis) that

(10) holds for each $x = (i, j) \in P$ satisfying $\mathbf{tilt} x < N$ and each $\ell \in \mathbb{N}$ satisfying $\ell - i - j + 1 \geq 0$ and $R^\ell f \neq \perp$.

We now fix an element $v = (i, j) \in P$ satisfying $\mathbf{tilt} v = N$ and an $\ell \in \mathbb{N}$ satisfying $\ell - i - j + 1 \geq 0$ and $R^\ell f \neq \perp$. Our goal is to prove that (10) holds for $x = v$. In other words, our goal is to prove that $v_\ell = a \cdot \overline{v_{\ell-i-j+1}^\sim} \cdot b$.

We have $N = \mathbf{tilt} v = i + 2j$ (since $v = (i, j)$). We are in one of the following six cases:

Case 1: We have $i = 1$.

Case 2: We have $j = 1$.

Case 3: We have $j = 2$ and $1 < i < p$.

Case 4: We have $j = 2$ and $i = p > 1$.

Case 5: We have $j > 2$ and $1 < i < p$.

Case 6: We have $j > 2$ and $i = p > 1$.

Let us first consider Case 1. In this case, we have $i = 1$. Thus, $v = (i, j) = (1, j)$ (since $i = 1$). The definition of v^\sim thus yields $v^\sim = (p+1-1, q+1-j) = (p, q+1-j)$. Also, $\ell - \underbrace{i}_{=1} - j + 1 = \ell - 1 - j + 1 = \ell - j$, so that $\ell - j = \ell - i - j + 1 \geq 0$. In other words, $\ell \geq j$. Hence, Lemma 10.2 yields

$$(1, j)_\ell = a \cdot \overline{(p, q+1-j)_{\ell-j}} \cdot b.$$

In view of $v = (1, j)$ and $v^\sim = (p, q + 1 - j)$ and $\ell - i - j + 1 = \ell - j$, we can rewrite this as $v_\ell = a \cdot \overline{v_{\ell-i-j+1}^\sim} \cdot b$. Thus, $v_\ell = a \cdot \overline{v_{\ell-i-j+1}^\sim} \cdot b$ is proved in Case 1.

Similarly (but using Lemma 10.1 instead of Lemma 10.2), we can obtain the same result (viz., $v_\ell = a \cdot \overline{v_{\ell-i-j+1}^\sim} \cdot b$) in Case 2.

Next, let us analyze the four remaining cases: Cases 3, 4, 5 and 6. The most complex of these four cases is Case 5, so it is this case that we start with.

In this case, we have $j > 2$ and $1 < i < p$. Recall that $v = (i, j)$. Define the four further pairs

$$\begin{aligned} m &:= (i, j - 1), & u &:= (i + 1, j - 1), \\ s &:= (i, j - 2), & t &:= (i - 1, j - 1). \end{aligned}$$

The conditions $j > 2$ and $1 < i < p$ entail that all these four pairs m, u, s and t belong to $[p] \times [q] = P$. Here is how the five elements v, m, u, s, t of P are aligned on the Hasse diagram of P :

$$\begin{array}{ccc} & u & v \\ & \diagdown & \diagup \\ & m & \\ & \diagup & \diagdown \\ s & & t \end{array} . \quad (46)$$

In particular, the two elements of \widehat{P} that cover m are u and v , whereas the two elements of \widehat{P} that are covered by m are s and t .

The map $P \rightarrow P, x \mapsto x^\sim$ (which can be visualized as “reflecting” each point in P around the center of the rectangle $[p] \times [q]$) “reverses” covering relations (i.e., if $x, y \in P$ satisfy $x \succ y$, then $x^\sim \prec y^\sim$). Hence, applying this map to the diagram (46) yields

$$\begin{array}{ccc} & t^\sim & s^\sim \\ & \diagdown & \diagup \\ & m^\sim & \\ & \diagup & \diagdown \\ v^\sim & & u^\sim \end{array} .$$

In particular, the two elements of \widehat{P} that are covered by m^\sim are u^\sim and v^\sim , whereas the two elements of \widehat{P} that cover m^\sim are s^\sim and t^\sim .

From $\ell - i - j + 1 \geq 0$, we obtain $\ell \geq \underbrace{i}_{>1} + \underbrace{j}_{>2} - 1 > 1 + 2 - 1 = 2$, so that $\ell \geq 2$.

Therefore, $\ell - 1 \in \mathbb{N}$ and $2 \leq \ell$.

Hence, from $R^\ell f \neq \perp$, we obtain $R^2 f \neq \perp$ (by Lemma 3.21). Therefore, Lemma 3.24 yields that a and b are invertible (since $a = f(0)$ and $b = f(1)$). Also, we have $R^{\ell-1} f \neq \perp$ (since $R(R^{\ell-1} f) = R^\ell f \neq \perp = R(\perp)$).

Set $k := i + j - 2$. Then, $k \geq 0$ (since $i \geq 1$ and $j \geq 1$), so that $k \in \mathbb{N}$.

Now, straightforward computations show that the four elements m, u, s and t of P satisfy

$$\mathbf{tilt} m < N, \quad \mathbf{tilt} u < N, \quad \mathbf{tilt} s < N, \quad \mathbf{tilt} t < N$$

(since $i + 2j = N$). Hence, using the induction hypothesis, it is easy to see that the five equalities

$$m_\ell = a \cdot \overline{m_{\ell-k}^\sim} \cdot b, \quad (47)$$

$$s_{\ell-1} = a \cdot \overline{s_{\ell-k}^\sim} \cdot b, \quad (48)$$

$$t_{\ell-1} = a \cdot \overline{t_{\ell-k}^\sim} \cdot b, \quad (49)$$

$$m_{\ell-1} = a \cdot \overline{m_{\ell-k-1}^\sim} \cdot b, \quad (50)$$

$$u_\ell = a \cdot \overline{u_{\ell-k-1}^\sim} \cdot b \quad (51)$$

hold²³.

We have $\ell - 1 \in \mathbb{N}$ and $R^\ell f \neq \perp$. Hence, the transition equation (12) (applied to m and $\ell - 1$ instead of v and ℓ) yields

$$m_\ell = \left(\sum_{x < m} x_{\ell-1} \right) \cdot \overline{m_{\ell-1}} \cdot \overline{\sum_{x > m} x_\ell}$$

(here we have renamed the summation indices u from (12) as x , since the letter u is already being used for something else in our current setting). Thus,

$$\begin{aligned} m_\ell &= \underbrace{\left(\sum_{x < m} x_{\ell-1} \right)}_{=s_{\ell-1}+t_{\ell-1}} \cdot \overline{m_{\ell-1}} \cdot \underbrace{\overline{\sum_{x > m} x_\ell}}_{=\overline{u_\ell+v_\ell}} \\ &\quad \text{(since the two elements of } \widehat{P} \text{ that are covered by } m \text{ are } s \text{ and } t) \quad \text{(since the two elements of } \widehat{P} \text{ that cover } m \text{ are } u \text{ and } v) \\ &= (s_{\ell-1} + t_{\ell-1}) \cdot \overline{m_{\ell-1}} \cdot \overline{u_\ell + v_\ell}. \end{aligned} \quad (52)$$

On the other hand, from $k = i + j - 2$, we obtain $\ell - k - 1 = \ell - i - j + 1 \geq 0$. Thus, $\ell - k - 1 \in \mathbb{N}$. Also, $\ell - \underbrace{k}_{\geq 0} \leq \ell$, so that $R^{\ell-k} f \neq \perp$ (by Lemma 3.21, since $R^\ell f \neq \perp$).

²³In more detail: The induction hypothesis tells us that ...

- ... we can apply (10) to m and $(i, j - 1)$ instead of x and (i, j) (since $m = (i, j - 1) \in P$ and tilt $m < N$ and $\ell \in \mathbb{N}$ and $\ell - i - \underbrace{(j - 1)}_{\leq j} + 1 \geq \ell - i - j + 1 \geq 0$ and $R^\ell f \neq \perp$). This yields (47)
(since an easy computation shows that $\ell - i - (j - 1) + 1 = \ell - k$).
- ... we can apply (10) to s and $(i, j - 2)$ and $\ell - 1$ instead of x and (i, j) and ℓ (the reader can easily verify that the requirements for this are satisfied). This yields (48) (since an easy computation shows that $(\ell - 1) - i - (j - 2) + 1 = \ell - k$).
- ... we can apply (10) to t and $(i - 1, j - 1)$ and $\ell - 1$ instead of x and (i, j) and ℓ . This yields (49).
- ... we can apply (10) to m and $(i, j - 1)$ and $\ell - 1$ instead of x and (i, j) and ℓ . This yields (50).
- ... we can apply (10) to u and $(i + 1, j - 1)$ instead of x and (i, j) . This yields (51).

Hence, the transition equation (12) (applied to m^\sim and $\ell - k - 1$ instead of v and ℓ) yields

$$\begin{aligned}
m_{\ell-k}^\sim &= \underbrace{\left(\sum_{x < m^\sim} x_{\ell-k-1} \right)}_{=u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim} \cdot \overline{m_{\ell-k-1}^\sim} \cdot \underbrace{\overline{\sum_{x > m^\sim} x_{\ell-k}}}_{=s_{\ell-k}^\sim + t_{\ell-k}^\sim} \\
&\quad \text{(since the two elements of } \widehat{P} \text{ that are covered by } m^\sim \text{ are } u^\sim \text{ and } v^\sim) \quad \text{(since the two elements of } \widehat{P} \text{ that cover } m^\sim \text{ are } s^\sim \text{ and } t^\sim) \\
&= (u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim) \cdot \overline{m_{\ell-k-1}^\sim} \cdot \overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}. \tag{53}
\end{aligned}$$

This entails that the elements $\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}$ and $m_{\ell-k-1}^\sim$ of \mathbb{K} are invertible (since their inverses appear on the right hand side of this equality). Hence, their product $(\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}) \cdot m_{\ell-k-1}^\sim$ is invertible as well.

Also, $\ell - k \geq 1$ (since $\ell - k - 1 \geq 0$) and $R^{\ell-k}f \neq \perp$. Hence, Lemma 7.1 (a) (applied to $\ell - k$ and m^\sim instead of ℓ and v) shows that $m_{\ell-k}^\sim$ is well-defined and invertible. Now, solving (53) for $u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim$, we obtain

$$u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim = m_{\ell-k}^\sim \cdot (\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}) \cdot m_{\ell-k-1}^\sim.$$

This shows that $u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim$ is invertible (since the three factors $m_{\ell-k}^\sim$ and $\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}$ and $m_{\ell-k-1}^\sim$ on the right hand side are invertible).

Taking reciprocals on both sides of (53), we obtain

$$\begin{aligned}
\overline{m_{\ell-k}^\sim} &= \overline{(u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim) \cdot m_{\ell-k-1}^\sim \cdot \overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}} \\
&= \overline{(\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}) \cdot m_{\ell-k-1}^\sim \cdot \overline{u_{\ell-k-1}^\sim + v_{\ell-k-1}^\sim}} \tag{54}
\end{aligned}$$

(by Proposition 2.3 (c)).

Comparing (52) with (47), we obtain

$$\begin{aligned}
a \cdot \overline{m_{\ell-k}^\sim} \cdot b &= \left(\underbrace{s_{\ell-1}}_{=a \cdot \overline{s_{\ell-k}^\sim} \cdot b} + \underbrace{t_{\ell-1}}_{=a \cdot \overline{t_{\ell-k}^\sim} \cdot b} \right) \cdot \underbrace{\overline{m_{\ell-1}^\sim}}_{=a \cdot \overline{m_{\ell-k-1}^\sim} \cdot b} \cdot \overline{u_\ell + v_\ell} \\
&= \underbrace{(a \cdot \overline{s_{\ell-k}^\sim} \cdot b + a \cdot \overline{t_{\ell-k}^\sim} \cdot b)}_{=a \cdot (\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}) \cdot b} \cdot \underbrace{\overline{a \cdot m_{\ell-k-1}^\sim \cdot b}}_{=b \cdot m_{\ell-k-1}^\sim \cdot a} \cdot \overline{u_\ell + v_\ell} \\
&\quad \text{(by Proposition 2.3 (c), since } a \text{ and } \overline{m_{\ell-k-1}^\sim} \text{ and } b \text{ are invertible)} \\
&= a \cdot (\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}) \cdot \underbrace{b \cdot \bar{b}}_{=1} \cdot m_{\ell-k-1}^\sim \cdot \bar{a} \cdot \overline{u_\ell + v_\ell} \\
&= a \cdot (\overline{s_{\ell-k}^\sim + t_{\ell-k}^\sim}) \cdot m_{\ell-k-1}^\sim \cdot \bar{a} \cdot \overline{u_\ell + v_\ell}.
\end{aligned}$$

Multiplying both sides of this equality by \bar{a} on the left and by \bar{b} on the right (this is

allowed, since a and b are invertible), we obtain

$$\begin{aligned} \overline{m_{\ell-k}^{\sim}} &= (\overline{s_{\ell-k}^{\sim}} + \overline{t_{\ell-k}^{\sim}}) \cdot m_{\ell-k-1}^{\sim} \cdot \underbrace{\overline{\overline{a \cdot \overline{u_{\ell}} + \overline{v_{\ell}} \cdot \overline{b}}}}_{= \overline{b \cdot (\overline{u_{\ell}} + \overline{v_{\ell}}) \cdot a}} \\ &\quad \text{(by Proposition 2.3 (c))} \\ &= (\overline{s_{\ell-k}^{\sim}} + \overline{t_{\ell-k}^{\sim}}) \cdot m_{\ell-k-1}^{\sim} \cdot \overline{b \cdot (\overline{u_{\ell}} + \overline{v_{\ell}}) \cdot a}. \end{aligned}$$

Comparing this with (54), we obtain

$$(\overline{s_{\ell-k}^{\sim}} + \overline{t_{\ell-k}^{\sim}}) \cdot m_{\ell-k-1}^{\sim} \cdot \overline{b \cdot (\overline{u_{\ell}} + \overline{v_{\ell}}) \cdot a} = (\overline{s_{\ell-k}^{\sim}} + \overline{t_{\ell-k}^{\sim}}) \cdot m_{\ell-k-1}^{\sim} \cdot \overline{u_{\ell-k-1}^{\sim} + v_{\ell-k-1}^{\sim}}.$$

Cancelling the $(\overline{s_{\ell-k}^{\sim}} + \overline{t_{\ell-k}^{\sim}}) \cdot m_{\ell-k-1}^{\sim}$ factors on the left of this equality (this is allowed, since $(\overline{s_{\ell-k}^{\sim}} + \overline{t_{\ell-k}^{\sim}}) \cdot m_{\ell-k-1}^{\sim}$ is invertible), we obtain

$$\overline{b \cdot (\overline{u_{\ell}} + \overline{v_{\ell}}) \cdot a} = \overline{u_{\ell-k-1}^{\sim} + v_{\ell-k-1}^{\sim}}.$$

Taking reciprocals on both sides, we find

$$b \cdot (\overline{u_{\ell}} + \overline{v_{\ell}}) \cdot a = u_{\ell-k-1}^{\sim} + v_{\ell-k-1}^{\sim}.$$

Expanding the left hand side by distributivity, we rewrite this as

$$b \cdot \overline{u_{\ell}} \cdot a + b \cdot \overline{v_{\ell}} \cdot a = u_{\ell-k-1}^{\sim} + v_{\ell-k-1}^{\sim}. \quad (55)$$

However, (51) yields

$$\overline{u_{\ell}} = \overline{a \cdot \overline{u_{\ell-k-1}^{\sim}} \cdot b} = \overline{b} \cdot u_{\ell-k-1}^{\sim} \cdot \overline{a} \quad \text{(by Proposition 2.3 (c))}.$$

Multiplying both sides by b from the left and by a from the right, we can transform this into

$$b \cdot \overline{u_{\ell}} \cdot a = u_{\ell-k-1}^{\sim}.$$

Subtracting this equality from (55), we obtain

$$b \cdot \overline{v_{\ell}} \cdot a = v_{\ell-k-1}^{\sim}. \quad (56)$$

This equality expresses $v_{\ell-k-1}^{\sim}$ as a product of three invertible elements (namely, b , $\overline{v_{\ell}}$ and a). Thus, $v_{\ell-k-1}^{\sim}$ is itself invertible.

Taking reciprocals on both sides of (56), we now obtain $\overline{b \cdot \overline{v_{\ell}} \cdot a} = \overline{v_{\ell-k-1}^{\sim}}$. Hence,

$$\overline{v_{\ell-k-1}^{\sim}} = \overline{b \cdot \overline{v_{\ell}} \cdot a} = \overline{a} \cdot v_{\ell} \cdot \overline{b} \quad \text{(by Proposition 2.3 (c))}.$$

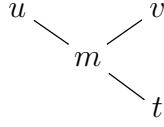
Solving this for v_{ℓ} , we obtain

$$v_{\ell} = a \cdot \overline{v_{\ell-k-1}^{\sim}} \cdot b = a \cdot \overline{v_{\ell-i-j+1}^{\sim}} \cdot b \quad \text{(since } \ell - k - 1 = \ell - i - j + 1 \text{)}.$$

Thus, $v_{\ell} = a \cdot \overline{v_{\ell-i-j+1}^{\sim}} \cdot b$ is proved in Case 5.

The arguments required to prove $v_{\ell} = a \cdot \overline{v_{\ell-i-j+1}^{\sim}} \cdot b$ in the Cases 3, 4 and 6 are similar to the one we have used in Case 5, but simpler:

- In Case 3, we have $s \notin P$. The “neighborhood” of m thus looks as follows:



(instead of looking as in (46)). This necessitates some changes to the proof; in particular, all addends that involve s or s^\sim in any way need to be removed, along with the equality (48).

- Case 6 is similar, but now we have $u \notin P$ instead. (Subtraction is no longer required in this case.)
- In Case 4, we have both $s \notin P$ and $u \notin P$.

Thus, we have proved the equality $v_\ell = a \cdot \overline{v_{\ell-i-j+1}^\sim} \cdot b$ in all six Cases 1, 2, 3, 4, 5 and 6. Hence, this equality always holds. In other words, (10) holds for $x = v$. This completes the induction step. Thus, (10) is proved by induction. In other words, Theorem 4.8 is proven. \square

As we have already seen (in Section 5), this entails that Theorem 4.7 is proven as well.

12. The case of a semiring

An attentive reader may have noticed that nowhere in the definitions of v -toggles and birational rowmotion do any subtraction sign appear. This means that all these definitions can be extended to the case when \mathbb{K} is not a ring but a *semiring*.

A *semiring* is a set \mathbb{K} equipped with a structure of an abelian semigroup $(\mathbb{K}, +)$ and the structure of a (not necessarily abelian) monoid $(\mathbb{K}, \cdot, 1)$ such that the distributive laws $(a + b)c = ac + bc$ and $a(b + c) = ab + ac$ are satisfied (where we use the shorthand notation xy for $x \cdot y$). Some standard concepts defined for rings can be straightforwardly generalized to semirings; in particular, any nonempty finite family $(a_i)_{i \in I}$ of elements of a semiring \mathbb{K} has a well-defined sum $\sum_{i \in I} a_i$. Definition 2.2, too, applies verbatim to the case when \mathbb{K} is a semiring instead of a ring. Thus, the definition of a v -toggle (Definition 3.10) and the definition of birational rowmotion (Definition 3.16) can be applied to a semiring \mathbb{K} as well. We thus can wonder:

Question 12.1. Do twisted periodicity (Theorem 4.7) and reciprocity (Theorem 4.8) still hold if \mathbb{K} is not a ring but merely a semiring?

If we assume that \mathbb{K} is commutative, then the answer to this question is positive, for fairly simple general reasons (see [GriRob16, Remark 10]). However, no such general reasoning helps for noncommutative \mathbb{K} . Indeed, there are subtraction-free identities involving inverses that hold for all rings but fail for some semirings. One example is the

identity $a \cdot \overline{a+b} \cdot b = b \cdot \overline{a+b} \cdot a$ from Proposition 2.4 (a): David Speyer has constructed an example of a semiring \mathbb{K} and two elements a and b of \mathbb{K} such that $a + b$ is invertible (actually, $a + b = 1$ in his example), but this identity does not hold. See [Speyer21] for details.

Of course, this does not mean that the answer to Question 12.1 is negative; we are, in fact, inclined to suspect that the question has a positive answer. Our proofs of Lemma 10.1 and Lemma 10.2 apply in the semiring setting (i.e., when \mathbb{K} is a semiring rather than a ring) without any need for changes; thus, Theorem 4.8 holds over any semiring \mathbb{K} at least in the case when one of i and j is 1. Unfortunately, subtraction is used in the proof of Theorem 4.8, and we have so far been unable to excise it from the argument. (With a bit of thought, we can convince ourselves that subtraction is actually unnecessary if $p = 2$ or $q = 2$, so the first interesting case is obtained for $P = [3] \times [3]$.)

13. Other posets: conjectures and results

We now proceed to discuss the behavior of R on some other families of posets P . We no longer use the notations introduced in Section 6.

13.1. The Δ and ∇ triangles

When $p = q$, the $p \times q$ -rectangle $[p] \times [q]$ becomes a square. By cutting this square in half along its horizontal axis, we obtain two triangles:

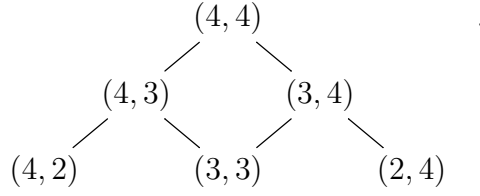
Definition 13.1. Let p be a positive integer. Define two subsets $\Delta(p)$ and $\nabla(p)$ of the $p \times p$ -rectangle $[p] \times [p]$ by

$$\begin{aligned}\Delta(p) &= \{(i, k) \in [p] \times [p] \mid i + k > p + 1\}; \\ \nabla(p) &= \{(i, k) \in [p] \times [p] \mid i + k < p + 1\}.\end{aligned}$$

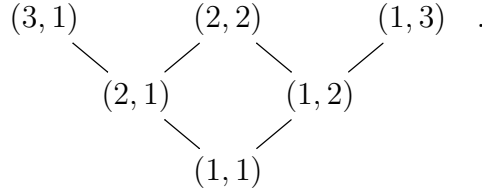
Each of these two subsets $\Delta(p)$ and $\nabla(p)$ inherits a poset structure from $[p] \times [p]$. In the following, we will consider $\Delta(p)$ and $\nabla(p)$ as posets using these structures.

The Hasse diagrams of these posets $\Delta(p)$ and $\nabla(p)$ look like triangles; if we draw $[p] \times [p]$ as agreed in Convention 4.4, then $\Delta(p)$ is the “upper half” of the square $[p] \times [p]$, whereas $\nabla(p)$ is the “lower half” of this square.

Example 13.2. Here is the Hasse diagram of the poset $\Delta(4)$:



Here, on the other hand, is the Hasse diagram of the poset $\nabla(4)$:



Note that $\Delta(p) = \emptyset$ when $p = 1$.

Computations with SageMath [S⁺09] for $p = 3$ have made us suspect a periodicity-like phenomenon similar to Theorem 4.7:

Conjecture 13.3 (periodicity conjecture for Δ -triangle). Let $p \geq 2$ be an integer. Assume that P is the poset $\Delta(p)$. Let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^p f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Let $x \in \widehat{P}$. We define an element $x' \in \widehat{P}$ as follows:

- If $x = 0$ or $x = 1$, then we set $x' := x$.
- Otherwise, we write x in the form $x = (i, j)$, and we set $x' := (j, i)$.

Then, a and b are invertible, and we have

$$(R^p f)(x) = a\bar{b} \cdot f(x') \cdot \bar{a}b.$$

Conjecture 13.4 (periodicity conjecture for ∇ -triangle). The same holds if $P = \nabla(p)$ instead of $P = \Delta(p)$.

If true, these two conjectures would generalize [GriRob15, Theorem 65], where \mathbb{K} is commutative.

13.2. The “right half” triangle

We can also cut the square $[p] \times [p]$ along its vertical axis:

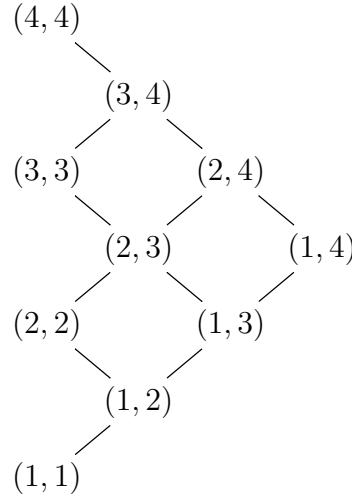
Definition 13.5. Let p be a positive integer. Define a subset $\text{Tria}(p)$ of the $p \times p$ -rectangle $[p] \times [p]$ by

$$\text{Tria}(p) := \{(i, k) \in [p] \times [p] \mid i \leq k\}.$$

This subset $\text{Tria}(p)$ inherits a poset structure from $[p] \times [p]$.

The Hasse diagram of this poset $\text{Tria}(p)$ has the shape of a triangle; if we draw $[p] \times [p]$ as agreed in Convention 4.4, then $\text{Tria}(p)$ is the “right half” of the square $[p] \times [p]$.

Example 13.6. Here is the Hasse diagram of the poset $\text{Tria}(4)$:



The inequality $i \leq k$ in Definition 13.5 could just as well be replaced by the reverse inequality $i \geq k$; the resulting poset would be isomorphic to $\text{Tria}(p)$. But we have to agree on something.

Now, we again suspect a periodicity-like phenomenon:

Conjecture 13.7 (periodicity conjecture for “right half” triangle). Let p be a positive integer. Assume that P is the poset $\text{Tria}(p)$. Let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^{2p}f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Let $x \in \widehat{P}$. Then, a and b are invertible, and we have

$$(R^{2p}f)(x) = a\bar{b} \cdot f(x) \cdot \bar{a}b.$$

If true, this conjecture would generalize [GriRob15, Theorem 58], where \mathbb{K} is commutative.

In a sense, we can “almost” prove Conjecture 13.7: Namely, the proof of its commutative case ([GriRob15, Theorem 58]) given in [GriRob15] can be adapted to the case of a general ring \mathbb{K} , as long as the number 2 is invertible in \mathbb{K} . The latter condition has all the earmarks of a technical assumption that should not matter for the validity of the result;

unfortunately, however, we are not aware of a rigorous argument that would allow us to dispose of such an assumption in the noncommutative case.

13.3. Trapezoids

Nathan Williams's conjecture [GriRob15, Conjecture 75], too, seems to extend to the noncommutative setting:

Conjecture 13.8 (periodicity conjecture for the trapezoid). Let p be an integer > 1 . Let $s \in \mathbb{N}$. Assume that P is the subposet

$$\{(i, k) \in [p] \times [p] \mid i + k > p + 1 \text{ and } i \leq k \text{ and } k \geq s\}$$

of $[p] \times [p]$. Let $f \in \mathbb{K}^{\widehat{P}}$ be a \mathbb{K} -labelling such that $R^p f \neq \perp$. Let $a = f(0)$ and $b = f(1)$. Let $x \in \widehat{P}$. Then, a and b are invertible, and we have

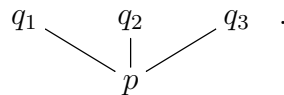
$$(R^p f)(x) = \bar{a}b \cdot f(x) \cdot \bar{a}b.$$

Again, this has been verified using SageMath for certain values of p and s and some randomly chosen \mathbb{K} -labellings with $\mathbb{K} = \mathbb{Q}^{3 \times 3}$. Even for commutative \mathbb{K} , a proof is yet to be found, although significant advances have been recently made on the lower (PL- and combinatorial) levels.

13.4. Ill-behaved posets

The above results and conjectures may suggest that every finite poset P for which birational rowmotion R has finite order when \mathbb{K} is commutative must also satisfy a similar (if slightly more complicated) property when \mathbb{K} is noncommutative. In particular, one might expect that if some positive integer m satisfies $R^m = \text{id}$ (as rational maps) for all fields \mathbb{K} , then $R^m f = f$ should also hold for all noncommutative rings \mathbb{K} and all \mathbb{K} -labellings $f \in \mathbb{K}^{\widehat{P}}$ that satisfy $f(0) = f(1) = 1$ (the latter condition ensures, e.g., that the $\bar{a}b$ and $\bar{a}b$ factors in Theorem 4.7 can be removed). However, this expectation is foiled by the following example:

Example 13.9. Let P be the four-element poset $\{p, q_1, q_2, q_3\}$ with order relation defined by setting $p < q_i$ for each $i \in \{1, 2, 3\}$. This poset has Hasse diagram



It is known (see [GriRob16, Example 18] or [GriRob16, Corollary 76]) that the birational rowmotion R of this poset P satisfies $R^6 = \text{id}$ (as rational maps) if \mathbb{K} is a field. In other words, if \mathbb{K} is a field, and if $f \in \mathbb{K}^{\widehat{P}}$ is a \mathbb{K} -labelling such that $R^6 f \neq \perp$, then

$R^6 f = f$. But nothing like this holds when \mathbb{K} is a noncommutative ring. For instance, if we let \mathbb{K} be the matrix ring $\mathbb{Q}^{2 \times 2}$, and if we define a \mathbb{K} -labelling $f \in \mathbb{K}^{\hat{P}}$ by

$$\begin{aligned} f(0) &= I_2 && \text{(the identity matrix in } \mathbb{K}\text{),} \\ f(1) &= I_2, & f(p) &= I_2, & f(q_1) &= I_2, \\ f(q_2) &= \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, & f(q_3) &= \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \end{aligned}$$

then $R^m f$ is distinct from f (and also distinct from \perp) for all positive integers m .

(See the detailed version of this article for a proof.)

Example 13.10. Let P be the four-element poset $\{p_1, p_2, q_1, q_2\}$ with order relation defined by setting $p_i < q_j$ for each i, j . It follows from [GriRob16, Proposition 74 (b) and Proposition 61] that the birational rowmotion R of this poset P satisfies $R^6 = \text{id}$ (as rational maps) if \mathbb{K} is a field. On the other hand, if \mathbb{K} is the matrix ring $\mathbb{Q}^{2 \times 2}$, then we can easily find a \mathbb{K} -labelling f of P such that $R^m f \neq f$ for all $1 \leq m \leq 10\,000$ (and probably for all positive m , but we have not verified this formally), despite $f(0)$ and $f(1)$ both being the identity matrix I_2 .

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